DBEst: Revisiting Approximate Query Processing Engines with Machine Learning Models

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Agenda

- Introduction
- AQP Engine
- DBEst Architecture
- DBEst Implementation
- Evaluation
- Related Work
- Conclusion & Future works

Introduction





What is the problem?

Selection Operators

Aggregate Functions

SELECT AF(y) FROM T WHERE x BETWEEN lb AND ub

Limitations?

AQP (Approximate Query Processing)



AQP State of Art

SAMPLING







AQP-Sampling

	Offline Sampling	Online Sampling	
Assumption	Workload is partially known.	No assumptions	
SpeedUp	High	Low	

AQP

Can we do AQP while ensuring: ?

YES!

- Small query execution time
- Small states (memory/storage)
- High accuracy
- Low money cost

DBEst

DEFINITION:

an AQP engine supporting the analytical needs, using prebuilt, a priori state



DBEst-What are the models?

TWO MODELS

D(x)

Density Estimator

R(x)

Regression Function

SELECT COUNT(sales_price,x) FROM store_sales WHERE (sales_price,x) BETWEEN Ib AND ub;

$$COUNT(y) \approx N \cdot \int_{lb}^{ub} D(x) dx$$

$$VARIANCE_x(x) = \mathbb{E} \left[x^2 \right] - \left[\mathbb{E} \left[x \right] \right]^2$$
$$= \frac{\int_{lb}^{ub} x^2 D(x) dx}{\int_{lb}^{ub} D(x) dx} - \left[\frac{\int_{lb}^{ub} x D(x) dx}{\int_{lb}^{ub} D(x) dx} \right]^2$$
$$STDDEV_x(x) = \sqrt{VARIANCE_x(x)}$$
$$= \sqrt{\frac{\int_{lb}^{ub} x^2 D(x) dx}{\int_{lb}^{ub} D(x) dx}} - \left[\frac{\int_{lb}^{ub} x D(x) dx}{\int_{lb}^{ub} D(x) dx} \right]^2$$

PERCENTILE $P(x < \alpha) = p$ $\int_{-\infty}^{\alpha} D(x) dx = p.$ F(x) = p**USE BISECTION!**

SELECT COUNT(x) FROM table WHERE x BETWEEN lb AND ub;

Problem: we need y to query!

USE $R_y(x)$

SELECT SUM(y) FROM

table WHERE x BETWEEN lb AND ub;

- - / >

SELECT AVG(sales_price) FROM store_sales WHERE sales_time BETWEEN lb AND ub;

$$AVG(y) = \mathbb{E}[y]$$

$$\approx \mathbb{E}[R(x)] \qquad E[f(x)] = \frac{\int_{lb}^{ub} f(x)D(x) \, dx}{\int_{-lb}^{ub} D(x) \, dx}$$

$$= \frac{\int_{lb}^{ub} D(x)R(x) \, dx}{\int_{lb}^{ub} D(x) \, dx}$$



EXTEND REGRESSION TO MULTIVAR:

SELECT SUM(sales_price) FROM store_sales WHERE sales_time BETWEEN Ib AND ub AND sold_time BETWEEN Ib AND ub; **R**₂(x, y)

SELECT SUM(z) FROM table WHERE x BETWEEN Ib AND ub AND yBETWEEN Ib AND ub;

SELECT z, AVG(y) FROM T WHERE x BETWEEN 1b AND ub GROUP BY z;

SOLUTION: Treat each z as having its own dataset to train model primitive on

DBEst-LIMITATIONS AND CHALLENGES

Models grow linearly with number of groups *increase query processing time? - parallelizable SOLUTION: create Model bundles to store model necessary for "High-cardinality" queries *still 10x as fast as sampling

DBEst Implementation

- 1. Sampling
- 2. Density Estimator
- 3. Regression Model Selection
- 4. Integral Evaluation
- 5. Parallel/Distributed Computation

Sampling

The paper mentions two sampling techniques



Density Estimator

- Kernel Estimator
 - High Accurate and Efficient
- Nearest neighbor method
- Orthogonal Series Estimators
- Histograms

Sample	1	2	3	4	5	6
Value	-2.1	- <mark>1.</mark> 3	- <mark>0.4</mark>	1.9	5.1	6.2



Regression Model Selection



Integral Evaluation

Interesting accuracy-efficiency trade-offs!



$$I_{w}[lb, ub]f = \int_{lb}^{ub} w(x)f(x)dx \qquad I_{w}[lb, ub]f \approx \sum_{k=1}^{n} w_{k}f(x_{k}) \qquad \{R_{n_{k}}, E_{n_{k}}\}, k = 1, 2, ..., N$$

w(.) is a weight function

x1,x2,...,xn are nodes, and w1, w2,, wn are weights

So what is happening here?

SELECT COUNT(pm25 real)

FROM mdl

WHERE PRES BETWEEN 1000 AND 1020;



Parallel/Distributed Computation

- Sampling -> easily parallelizable
 - Different nodes storing dataset partitions
- Model Training -> easily parallelizable
 - GROUP BY queries
- Query Processing -> easily parallelizable
 - Additional nodes/cores



Performance Evaluation

- 1. Experimental Setup: Ubuntu 18.04 with Intel Xeon X5650 12-core CPU, 64GB RAM and 4TB of SSD
- 2. Datasets: TPC-DS, Combined Cycle Power Plant(CCPP), Beijing PM2.5(UCI-ML)
- 3. Query Types:
 - Synthetic queries: w/ 0.1%, 1%, and 10% query range
 - Complex TPC-DS queries
- 4. Comparison:
 - w/ VerdictDB, Blink DB and MonetDB for error
 - \circ w/VerdictDB for time
- 5. Additional:
 - VerdictDB uses 12 cores while DBEst runs on 1 core(Multi-threaded DBEst is also evaluated)
 - Performance of joins and group by

Sensitivity Analysis (Query Range & Sample Size)



Figure 5: Influence of Query Range on Relative Error

- Dataset:TPC-DS
- Sample size: 100k rows
- Query Range: 0.1%, 1%, 10%



Figure 2: Influence of Sample Size on Relative Error

- Dataset:TPC-DS
- Query range: 1%
- Sample size: 10k, 100k, 1m, 5m



Figure 4: DBEst vs VerdictDB Overheads

• 1 to 2 orders of magnitude less than VerdictDB's

Performance comparison TPC-DS & CCPP dataset



TPC-DS







Figure 9: Response Time for CCPP Dataset

Figure 10: Relative Error: DBEst vs VerdictDB Figure 11: Response Time: DBEst vs VerdictDB

- Query range: 0.1%
- Sample size: 10k, 100k
- Response time is small

- Sample size: 10k, 100k
- Reminder: VerdictDB uses 12 cores, while DBEst uses 1 thread

Performance comparison GROUP BY



- 90 queries, 57 groups
- Sample Size: 10k



• VerdictDB has no benefit from Parallel version

Performance comparison Join



Figure 20: Join Accuracy Comparison



Figure 21: Join Performance Comparison

• Sample size: 10k, 100k, 1m for DBEst; 10m for VerdictDB

Limitation

- **GROUP BY queries:** As the number of group increase
 - Models ↑
 - \circ Training time \uparrow
 - \circ Query Response Time \uparrow
 - \circ Space overheads \uparrow
- Small groups:
 - Building Models is an overkill
- No error guarantees



Conclusion & Future Works

• DBEst:

- Smaller query response time
- Higher accuracy
- Smaller space-time overheads
- scalability
- All comes with low money cost!!
- Future:
 - Offering better efficiency-overheads-accuracy trade-offs(especially Joins queries)
 - categorical attributes
 - parallel/distributed DBEst