BOSTON UNIVERSITY

CS 561: Data Systems Architectures

class 7

Row-stores vs. Column-stores

Dr. Subhadeep Sarkar

https://bu-disc.github.io/CS561/

with slides based on Dan Abadi's

A few reminders

A) Project 1 is out – start working on it now!

B) Submissions **due on 02/20** in groups of 2-3.

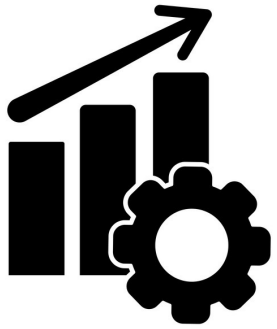
C) Class projects to be announced next week (form groups).

D) First review due on **02/14**.

E) First student presentation **02/14** (discussion + Q&A).

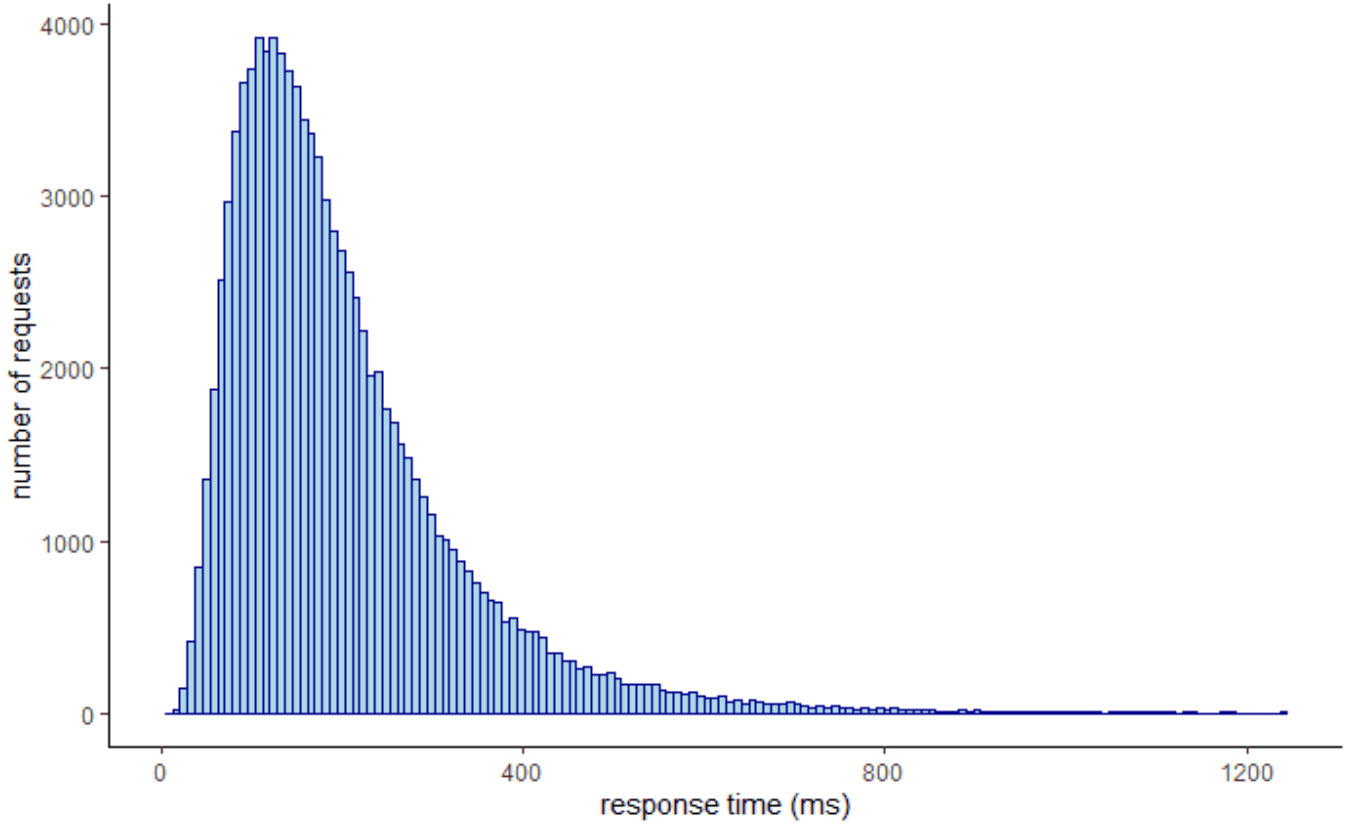
Performance Metrics

why care about both?



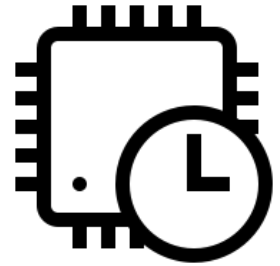
throughput

latency

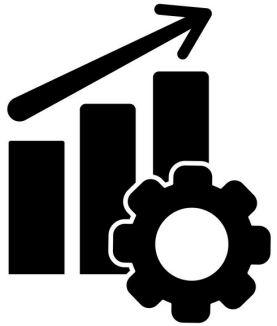


Performance Metrics

I/O

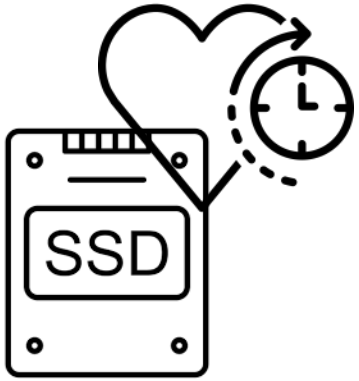


CPU

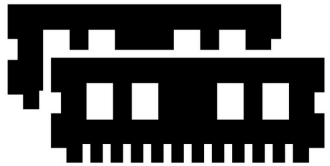


throughput

latency



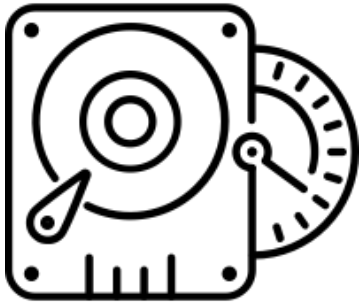
device lifetime



memory
miss/hit

cache
miss/hit

stalls



bandwidth
utilization

Row-stores vs. Col-Stores: How Different Are They Really?

Are column-stores really novel?

If we profile their performance, what is the breakdown? Why?

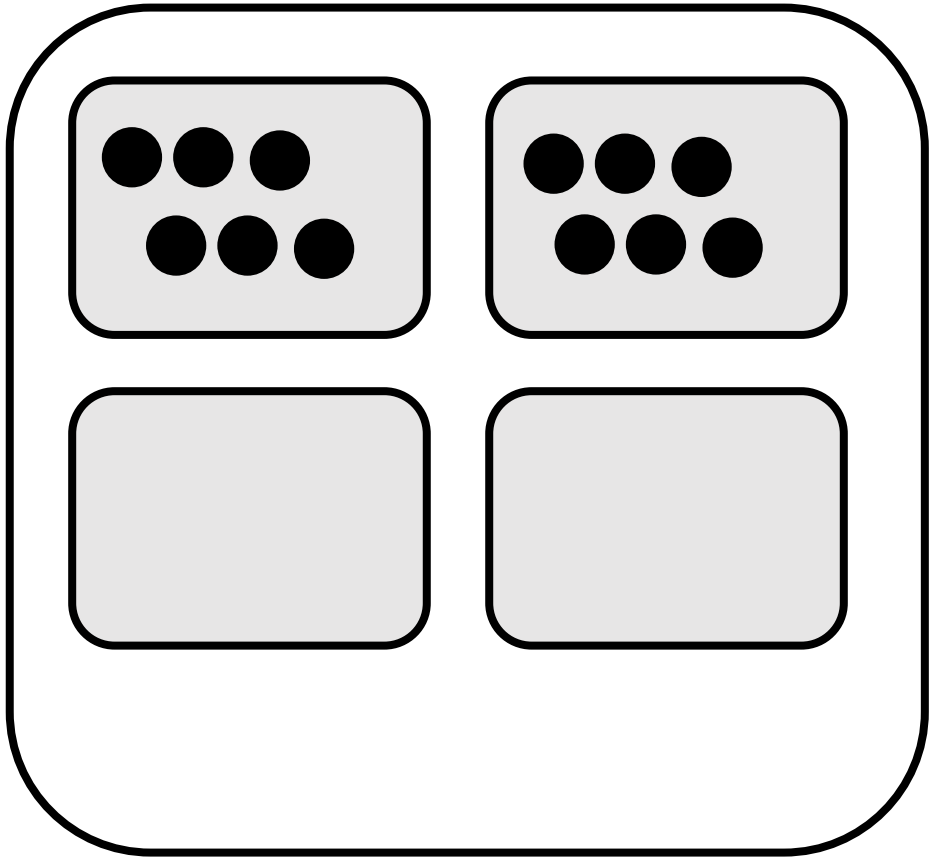
The paper tries to clarify which part of the “column stores” hype was marketing and which was fundamental

Row-Stores

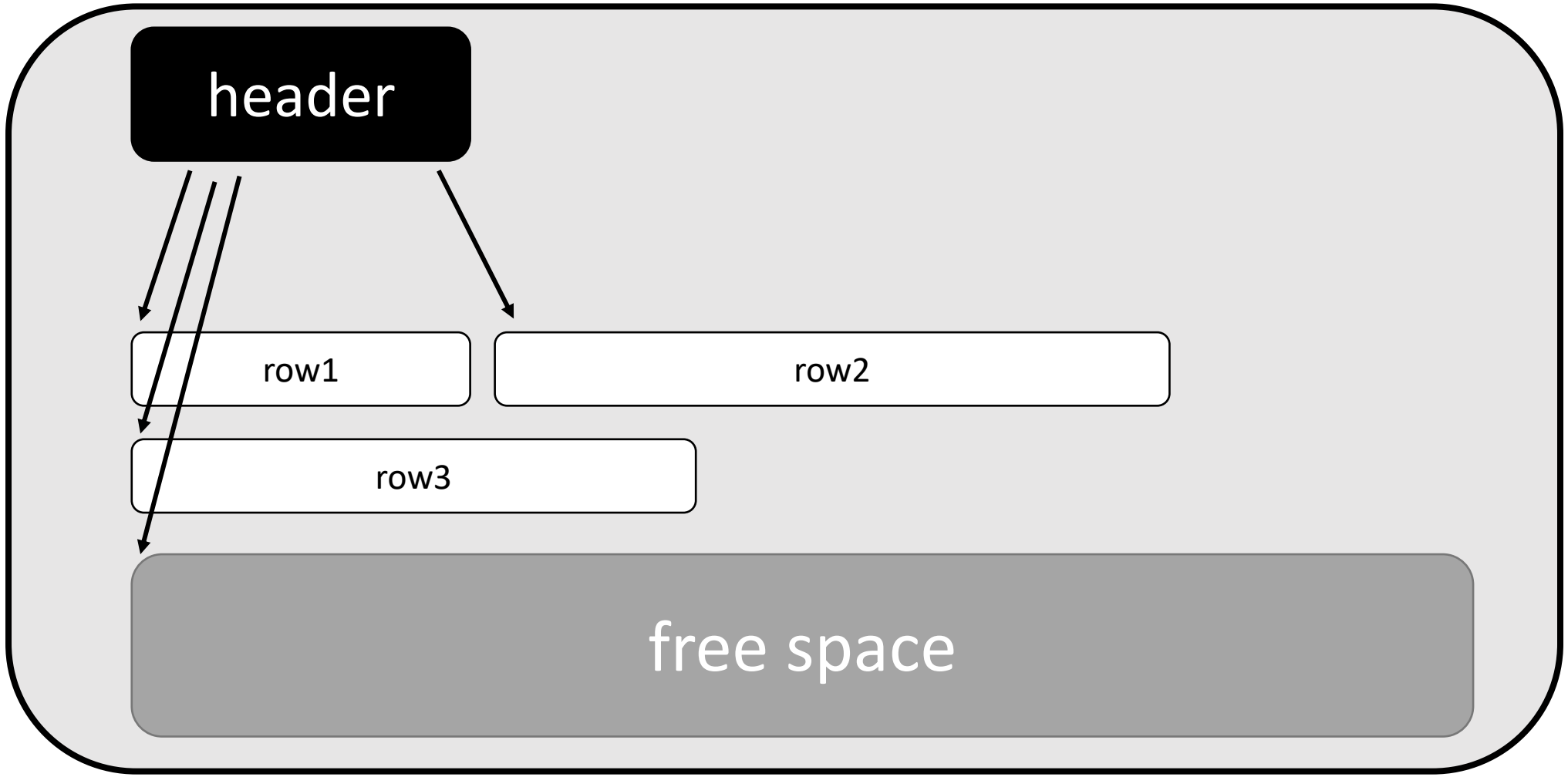
*Student (**sid**: string, **name**: string, **login**: string, **year_birth**: integer, **gpa**: real)*

student

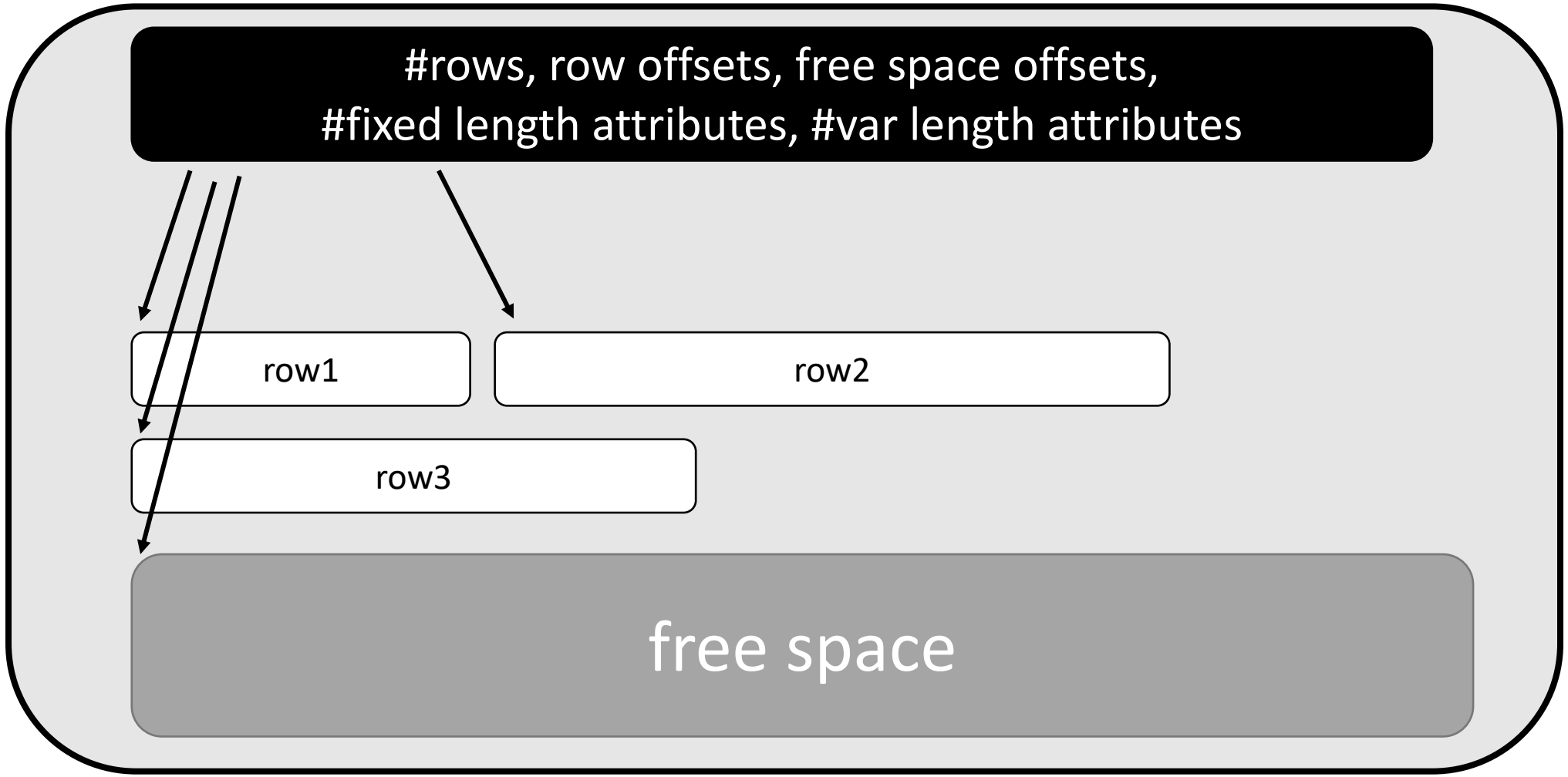
(sid1, name1, login1, year1, gpa1)
(sid2, name2, login2, year2, gpa2)
(sid3, name3, login3, year3, gpa3)
(sid4, name4, login4, year4, gpa4)
(sid5, name5, login5, year5, gpa5)
(sid6, name6, login6, year6, gpa6)
(sid7, name7, login7, year7, gpa7)
(sid8, name8, login8, year8, gpa8)
(sid9, name9, login9, year9, gpa9)



Row-Stores: slotted page

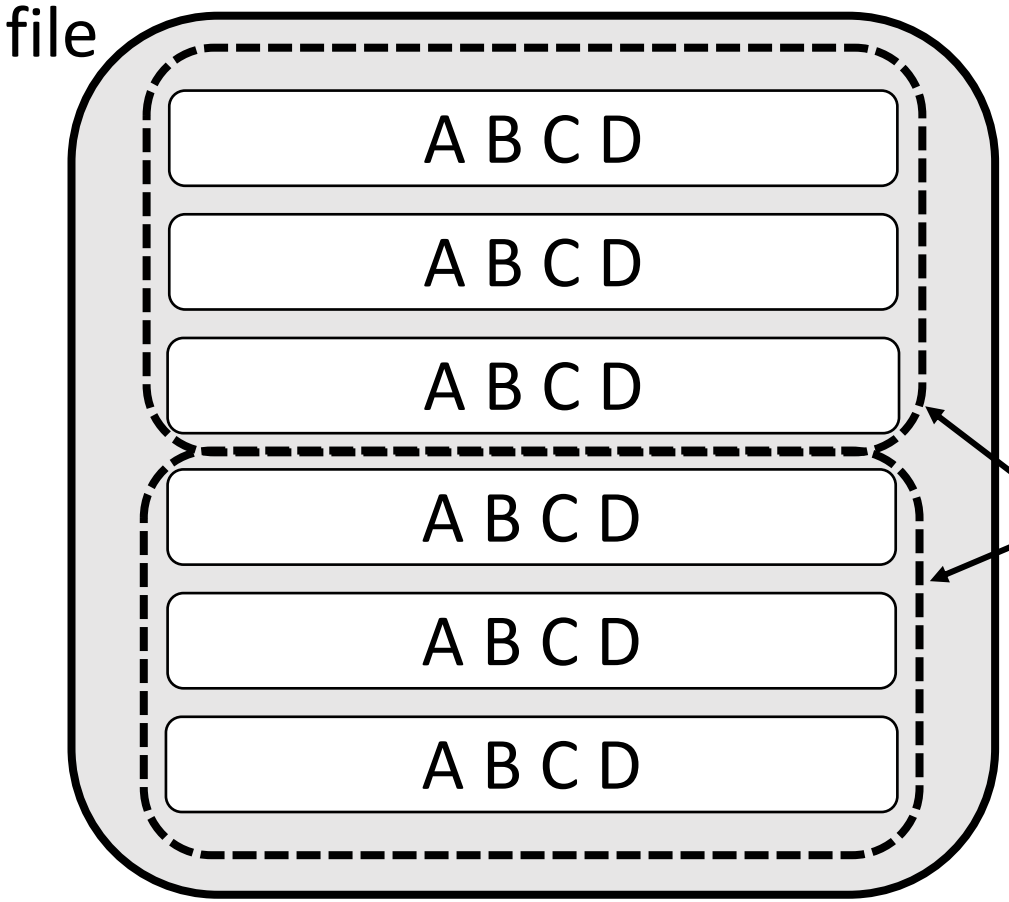


Row-Stores: slotted page



Row-Stores

pros and cons?



✅ Easy to add a new record

❌ Might access unnecessary data

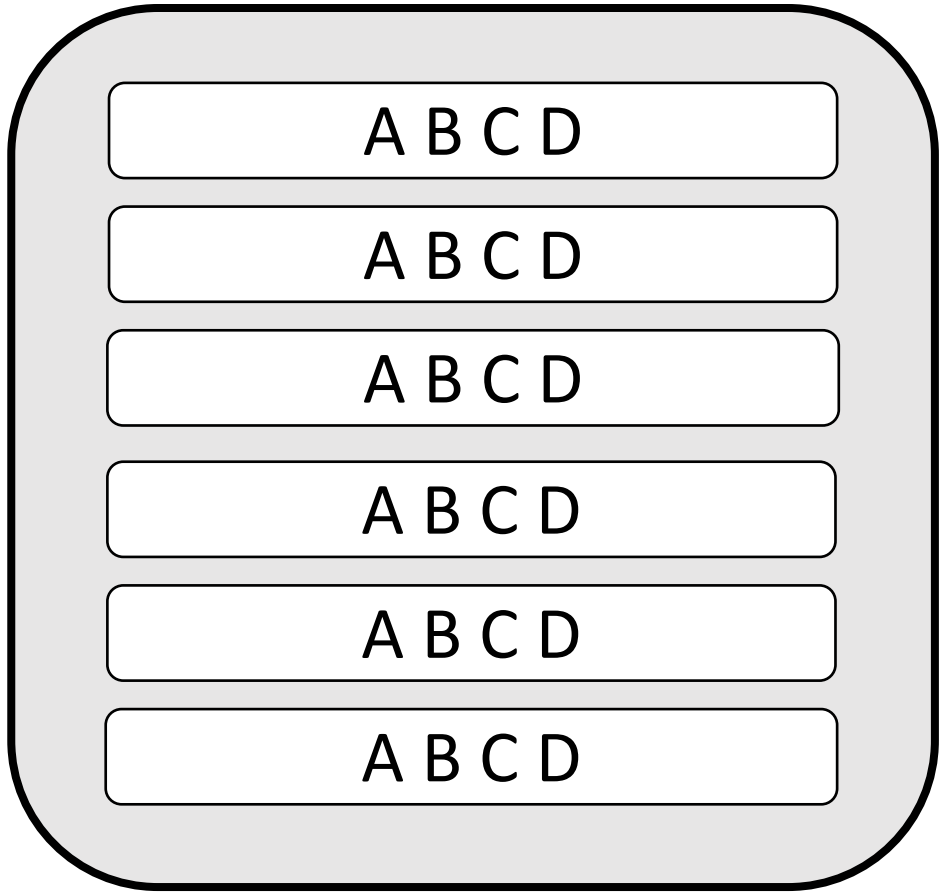
each page contains **entire** rows (all their columns)

pages

rows are **contiguous**
(with possible free space at the end)

Row-stores: query processing

select max(B) **from** R **where** A>5 **and** C<10



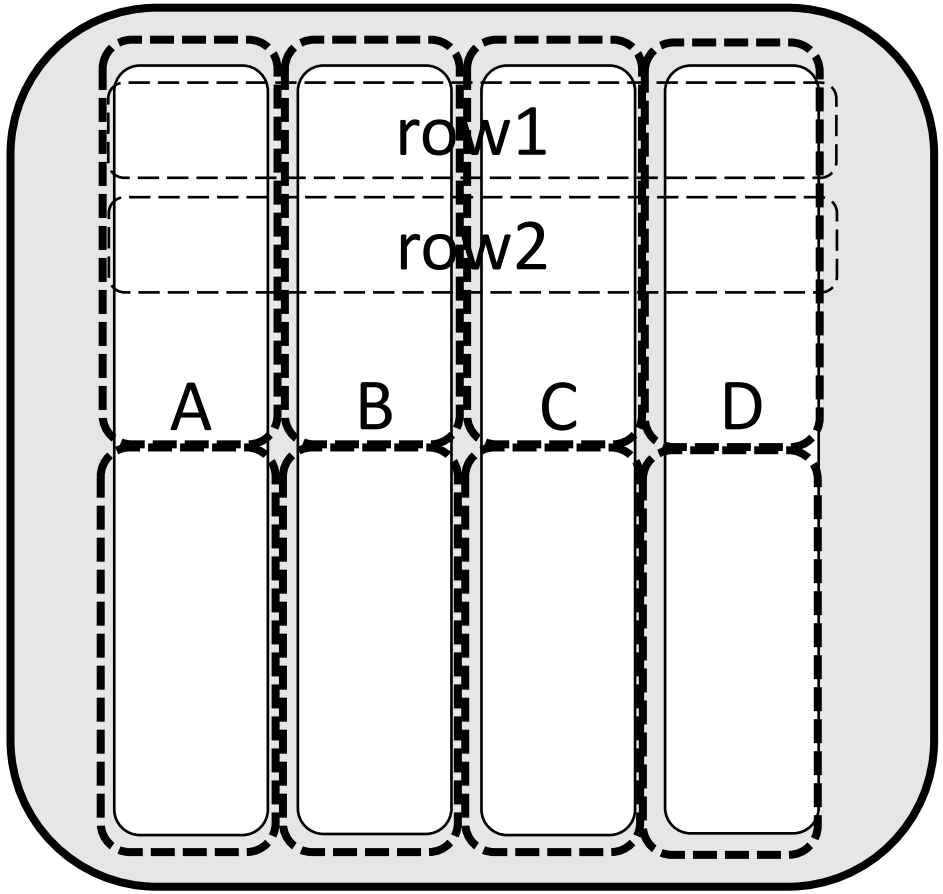
A B C D

one row at a time

what's that?

Early materialization

Column-Stores



pros and cons?

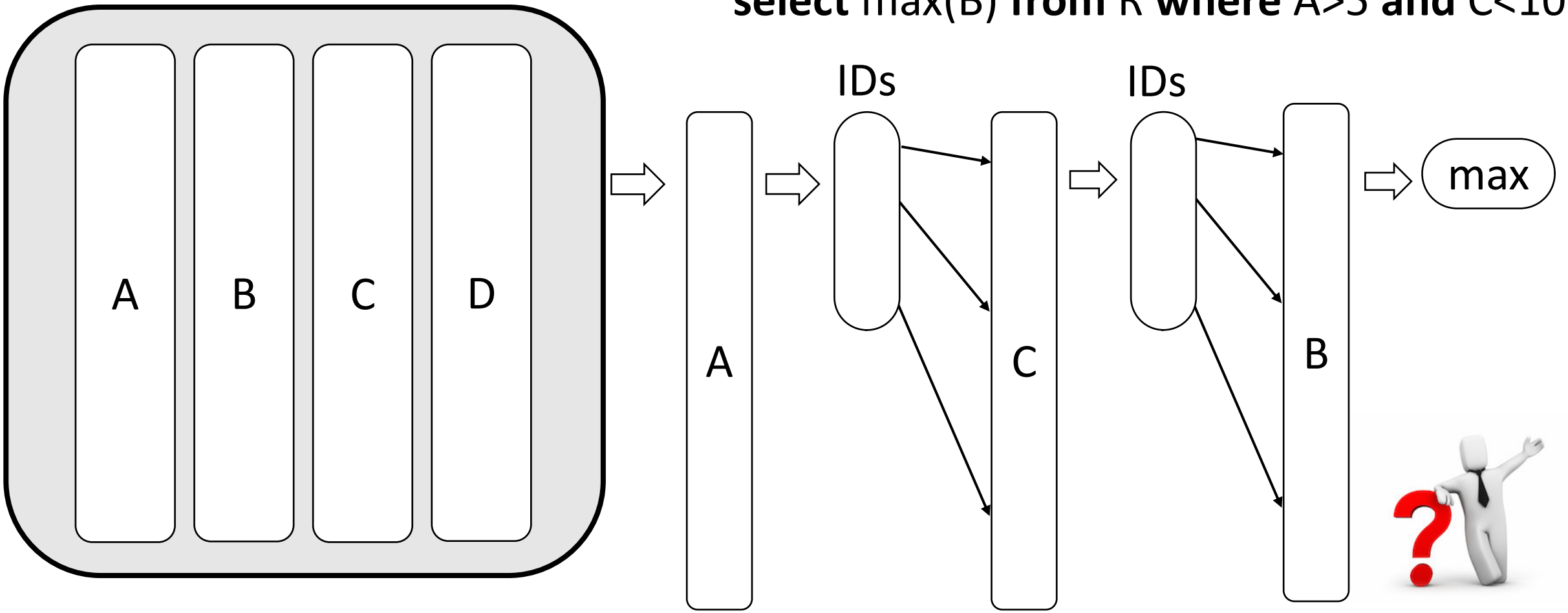
✅ Read only relevant data

❌ Tuple writes require multiple accesses

each page contains **columns!**

Column-stores: query processing

select max(B) **from** R **where** A>5 **and** C<10



why so?

Late materialization

Let's revisit the main question of the paper

What's the goal of the paper?

Prior to this paper there several studies showing

column-stores outperforming row-stores (~5x better performance in TPCH)

Compare row-stores and column-stores

Key question:

(a) are the benefits **inherent to the new column-store design**, or

(b) a **row-store with a "more columnar"** physical design can **achieve the same**?

In other words: ***can you "simulate a col-store in a row-store?"***

Paper's Methodology

Compare row-store vs. row-store and col-store vs. col-store.

How?

1. Simulate a column-store inside a row-store
2. Remove col-store features one-by-one

How to simulate a col-store with a row-store?

Vertical Partitioning

> "physically partition the data per column"

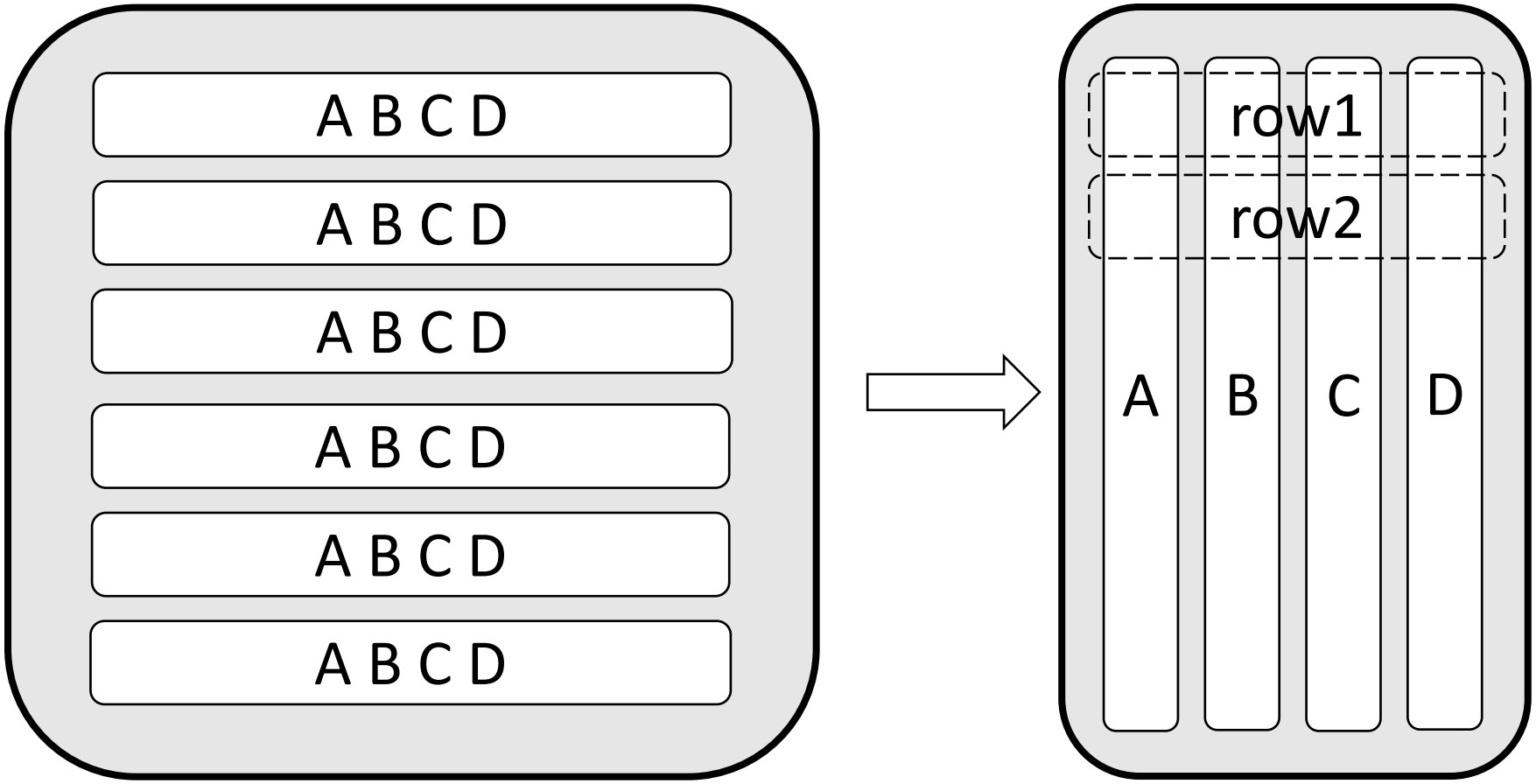
Index-only Plans

> "use only indexes in query plans that contain only relevant columns"

Materialized Views

> "temporary tables that contain exactly the answer to a query"

Vertical Partitioning



Problem?

Details on Vertical Partitioning

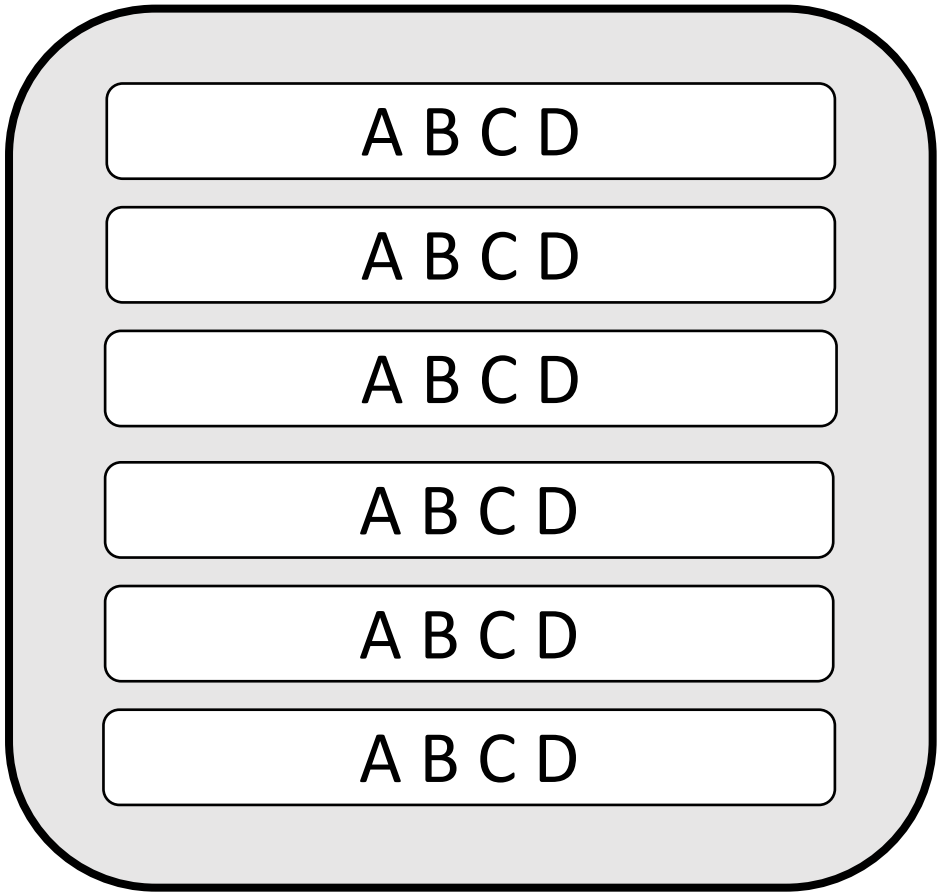
TID	Column Data
1	
2	
3	

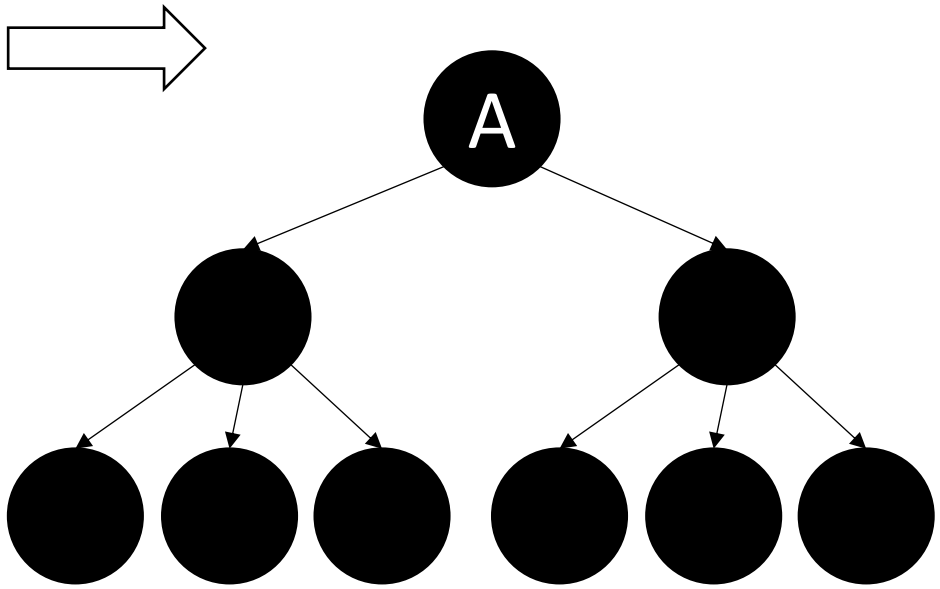
TID	Column Data
1	
2	
3	

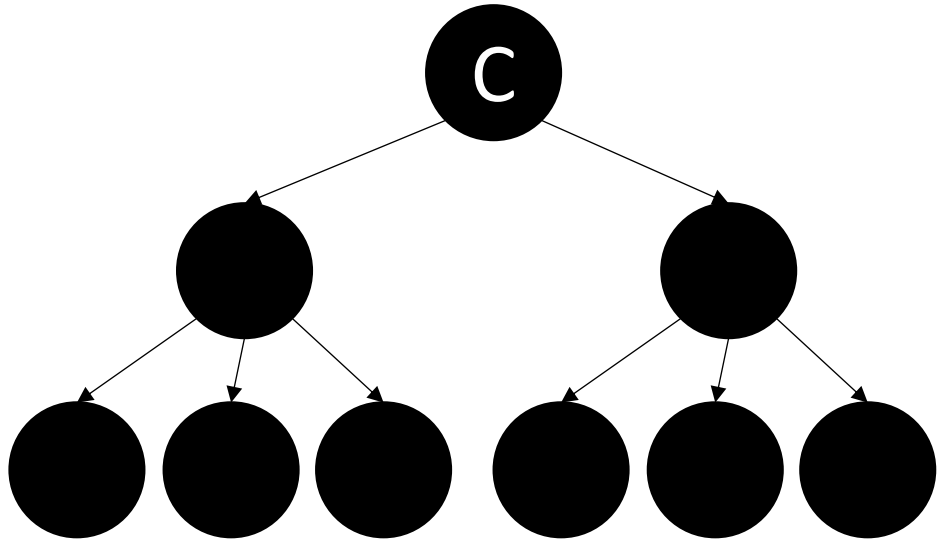
Tuple Header	TID	Column Data
	1	
	2	
	3	

Note that a "real column-store" would only store the raw values as an array.

Index-only plans

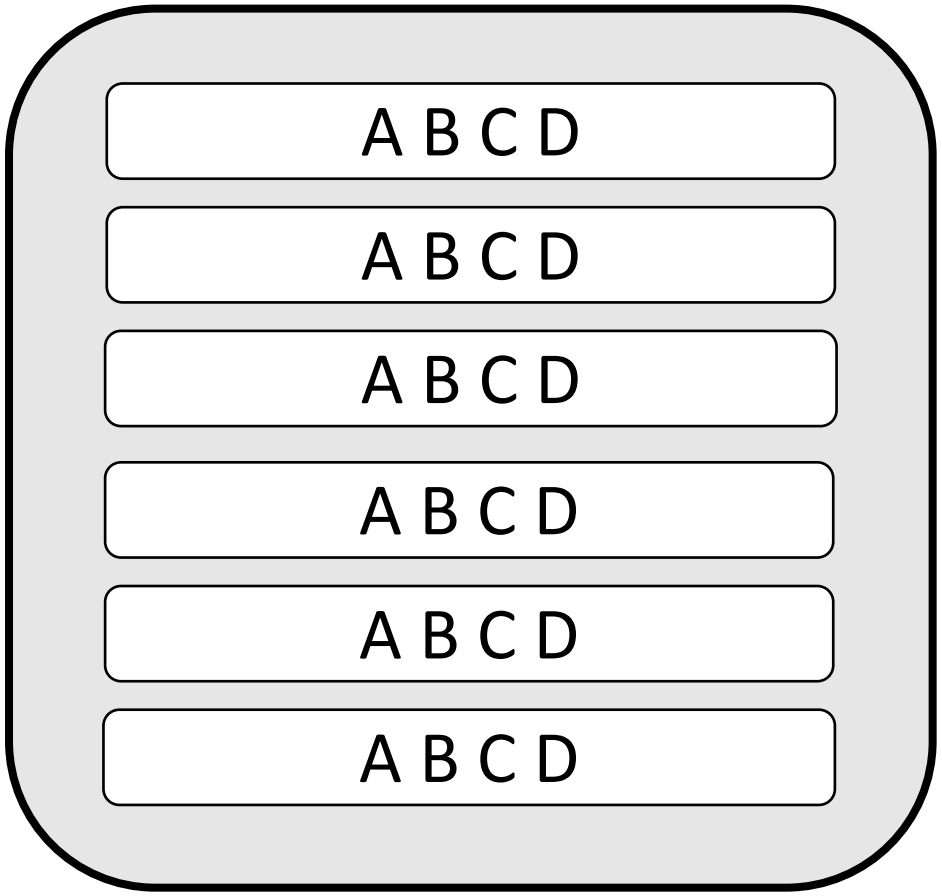


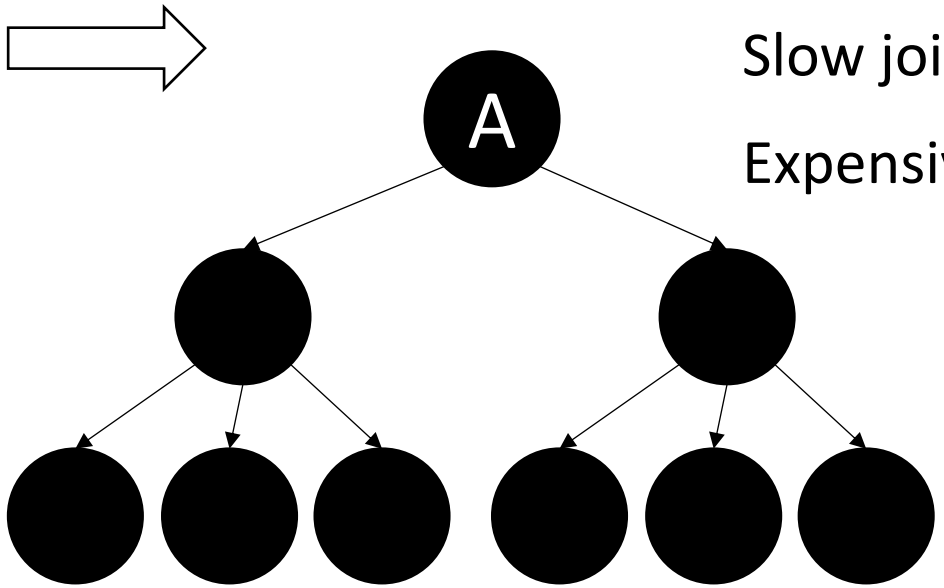


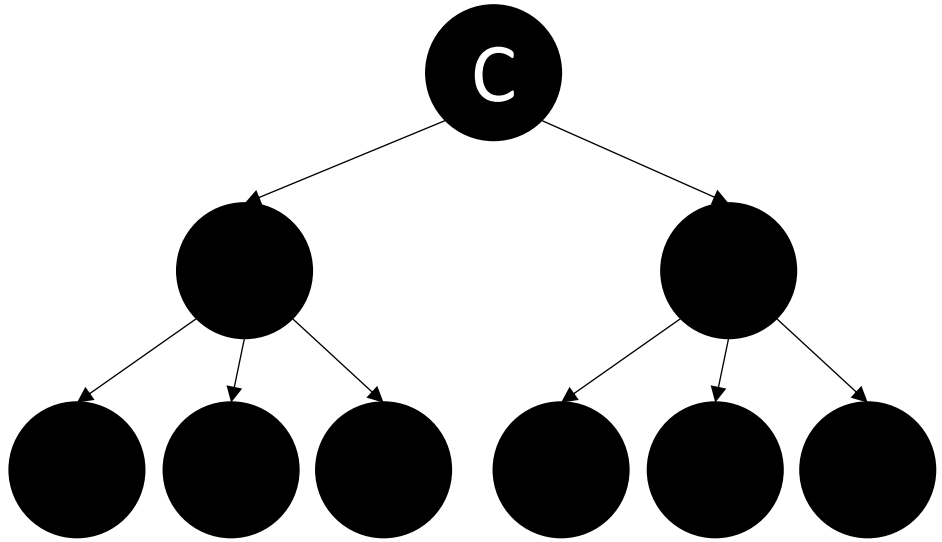


Problem?

Index-only plans



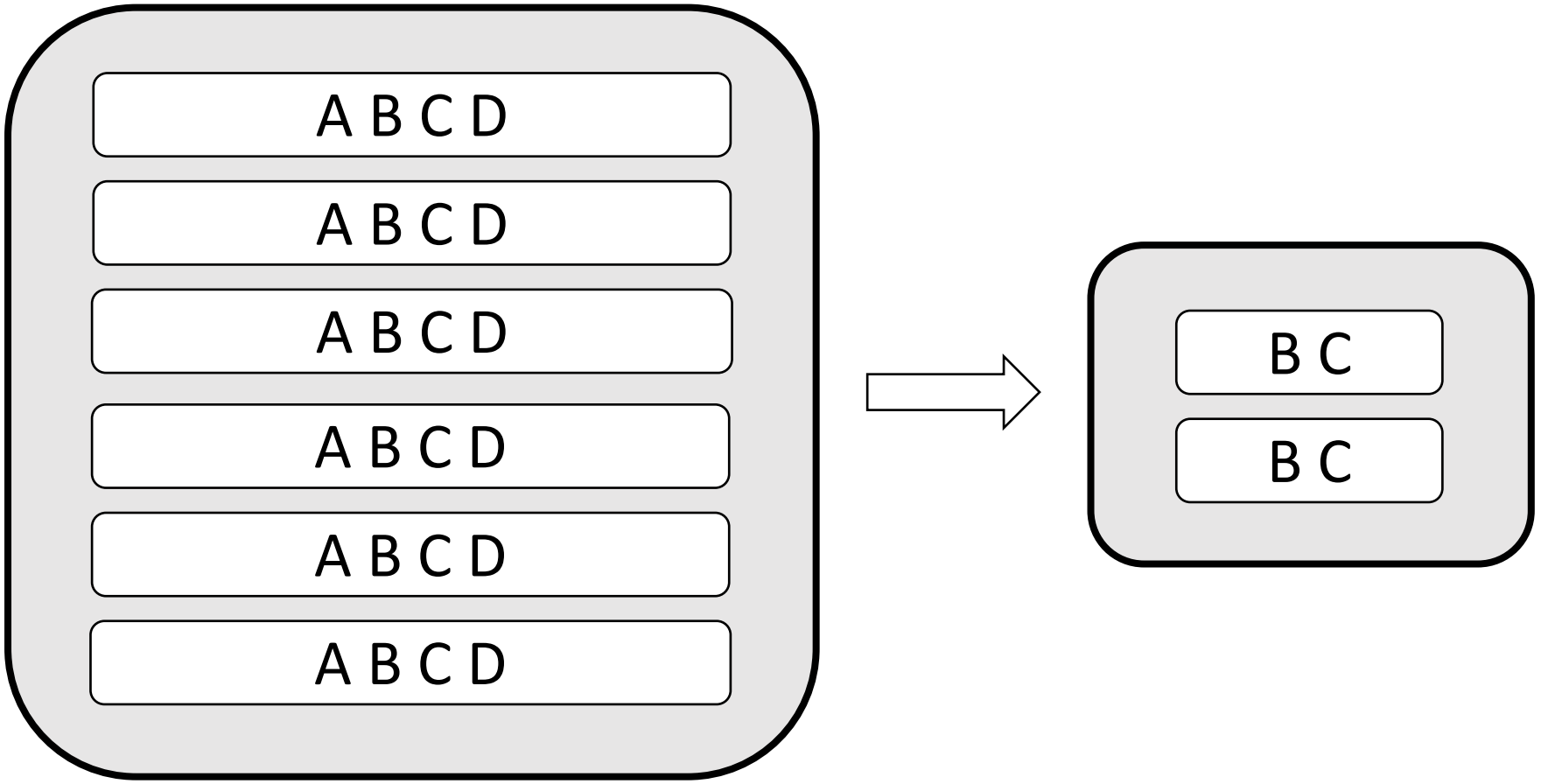




Slow joins

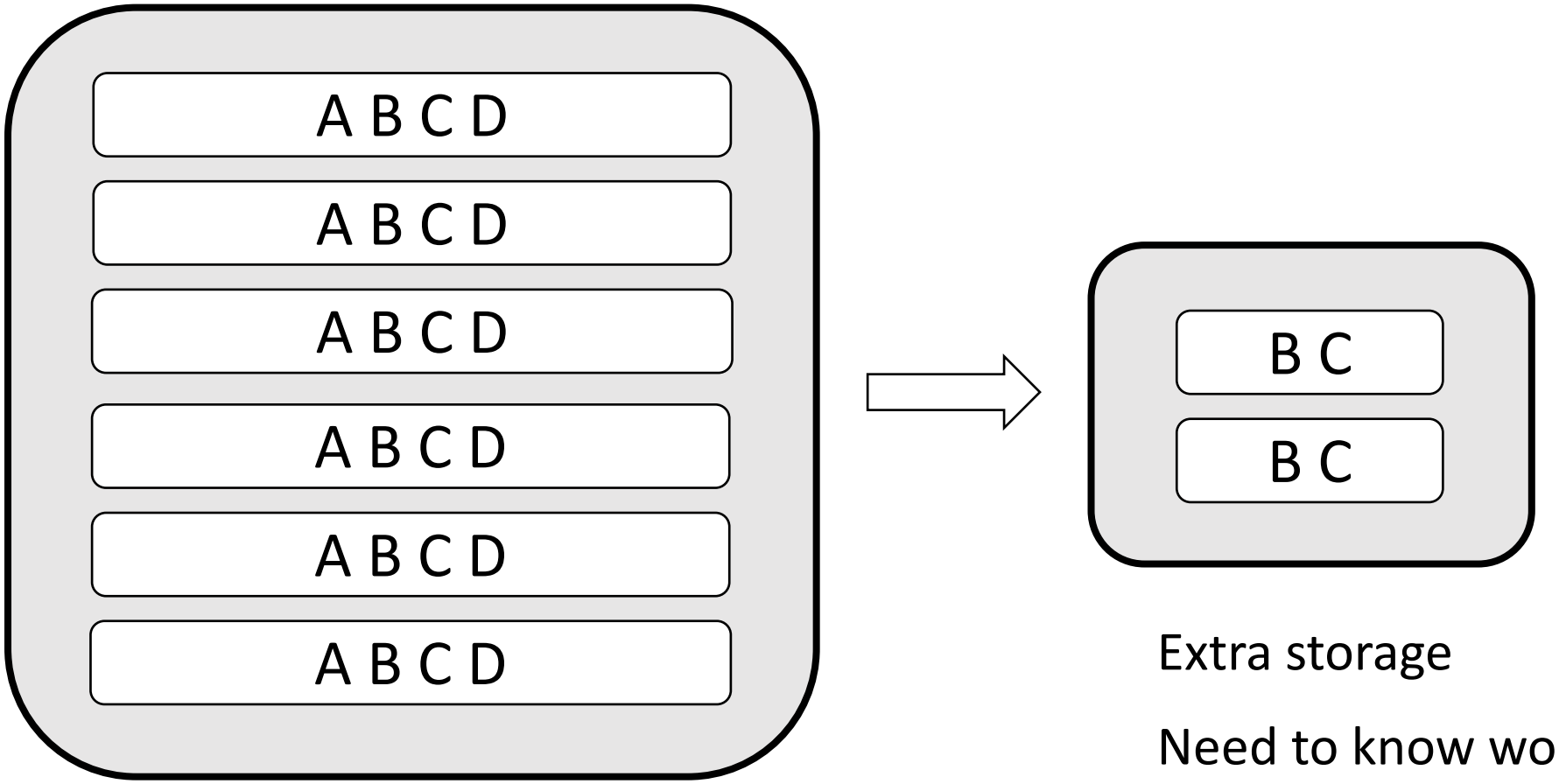
Expensive random access

Materialized Views



Problem?

Materialized Views



Extra storage

Need to know workload apriori

State-of-the-art Col-Store features

Late Materialization

"stich the column together as late as possible"

Block iteration

"execute the same columnar operation over a block of values"

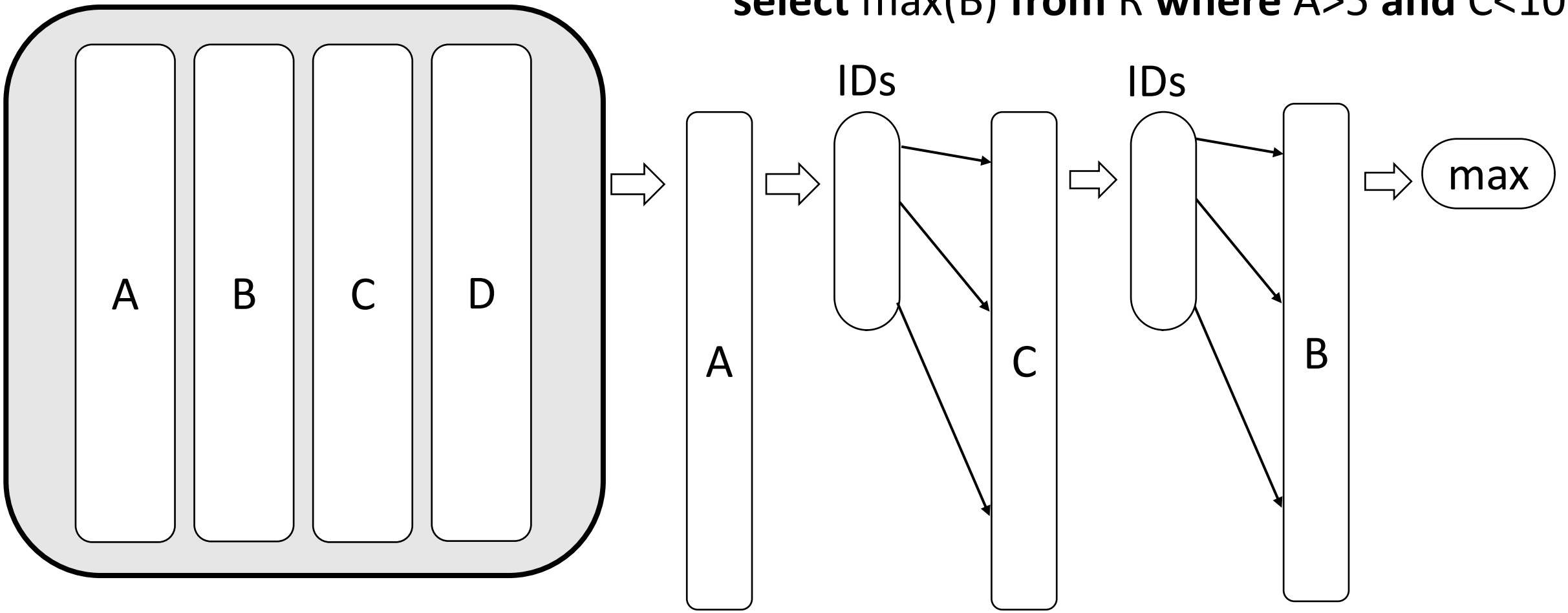
Compression

"column-specific compression, due to the nature of data"

Invisible joins

Late Materialization

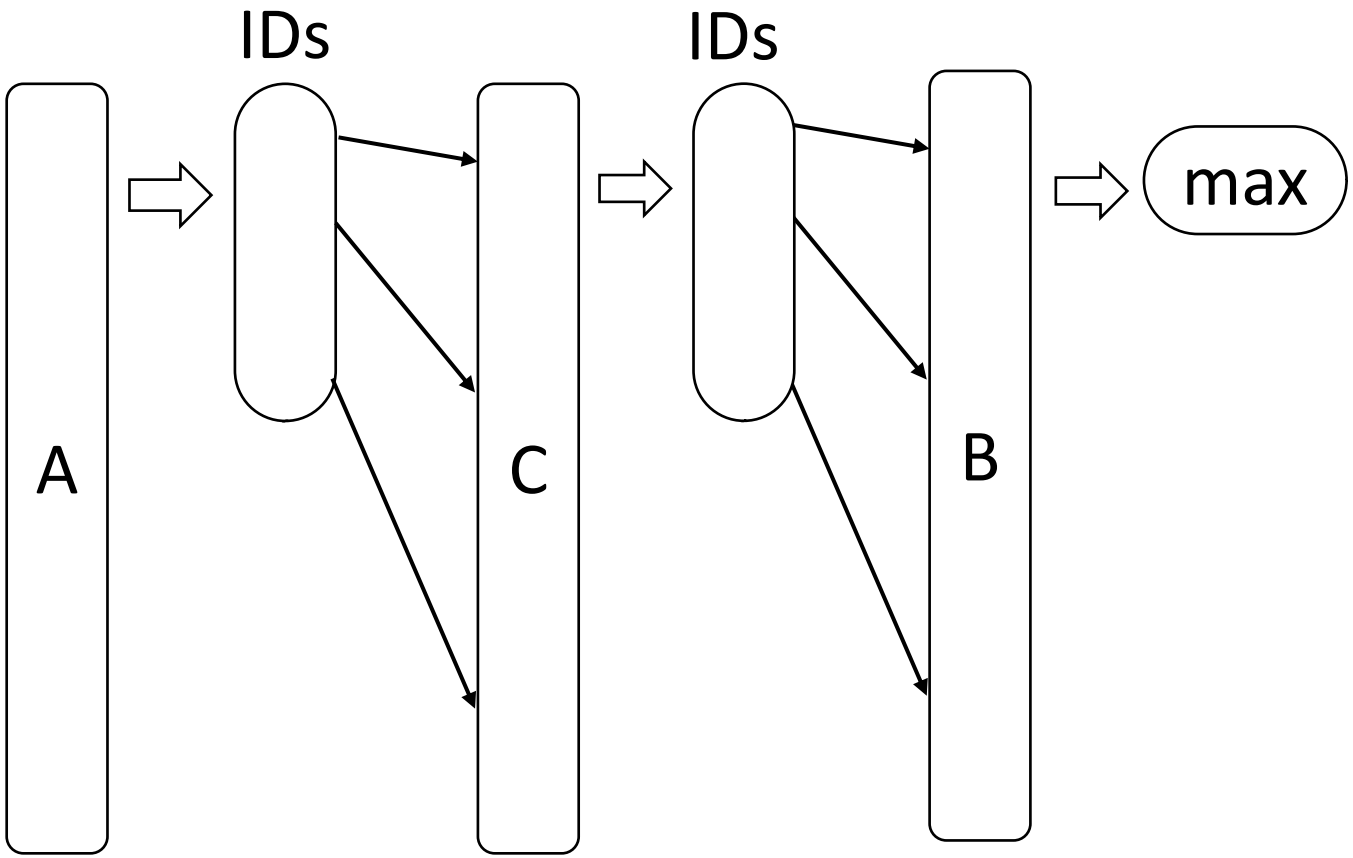
select max(B) **from** R **where** A>5 **and** C<10



"the full tuple (or the necessary subset) is not materialized until it is needed"

“Column-at-a-time”

select max(B) **from** R **where** A>5 **and** C<10



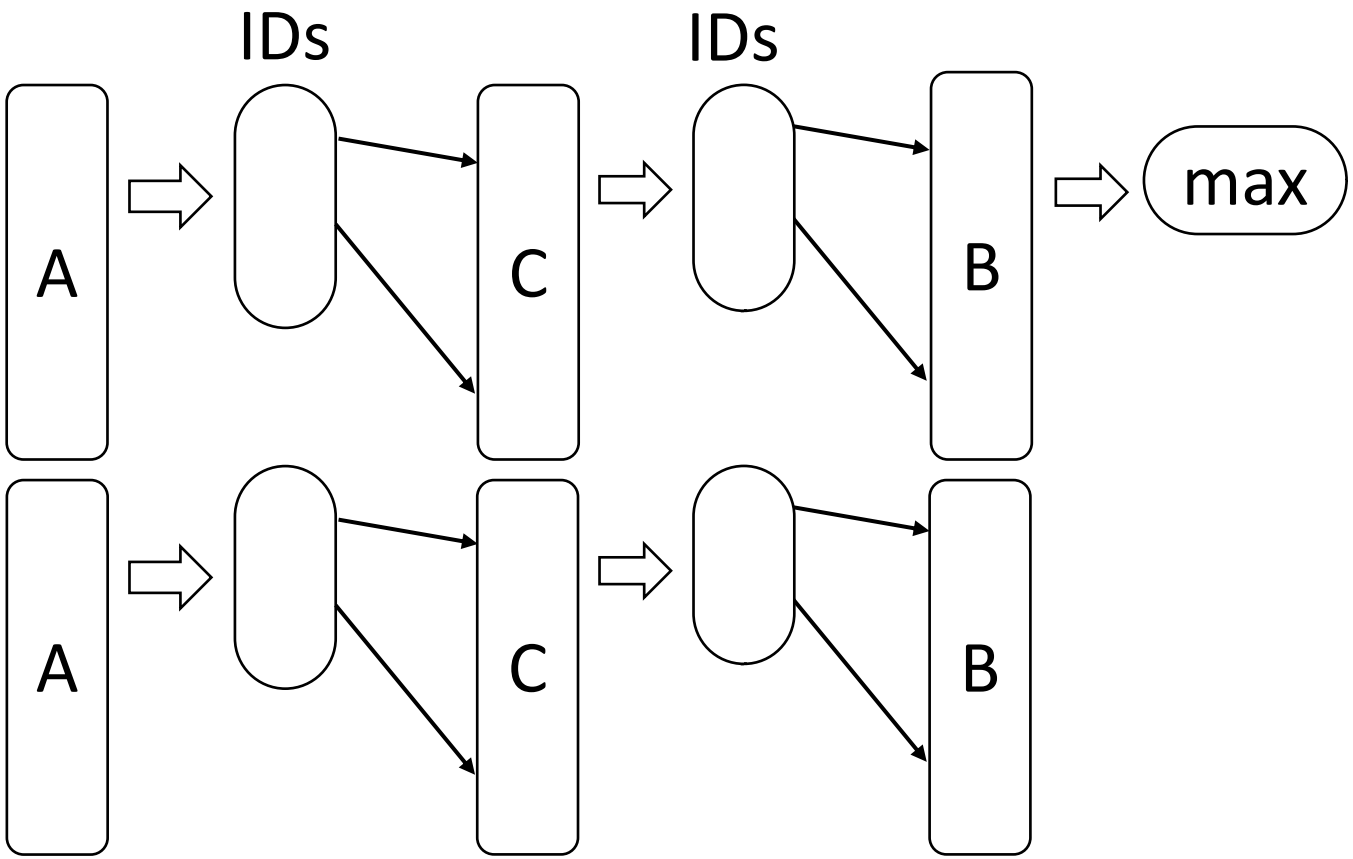
whole column?

column at a time

block/vector at a time

Block Iteration

select max(B) **from** R **where** A>5 **and** C<10



whole column?

column at a time

block/vector at a time

What is easier to compress?

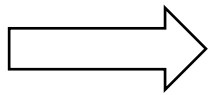
#1, John, 2/4/88, Boston			
#2, Joe, 2/1/87, New York			
#3, Lina, 7/7/93, Boston			
#4, Anna, 4/1/92, Chicago			
#5, Tim, 3/9/91, Seattle			
#6, Rose, 9/3/96, Boston			

#1	John	2/4/88	Boston
#2	Joe	2/1/87	New York
#3	Lina	7/7/93	Boston
#4	Anna	4/1/92	Chicago
#5	Tim	3/9/91	Seattle
#6	Rose	9/3/96	Boston

exploit patterns, duplicates, small differences

Compression

Q1
Q1
Q1
...
Q2
Q2
...



Run-length Encoding

Q1, 1, 300
Q2, 301, 500
Q3, 501, 550
Q4, 551, 800

Alternative: Dictionary Compression

- Replace variable size with minimal fixed length e.g., integer

Benefits of col-store compression

Reduces I/O

Can operate directly on compressed data

How?

Are the same benefits applicable for row-store compression?

Reduces I/O → yes, but with lower ratio

No! Requires decompression before processing

Invisible Joins

Benchmarking

When comparing database systems we need a common “language”

Benchmarks from the ***Transaction Performance Council***

TPC-B, TPC-C, TPC-H, TPC-DS etc

Also, a benchmark for data warehousing:

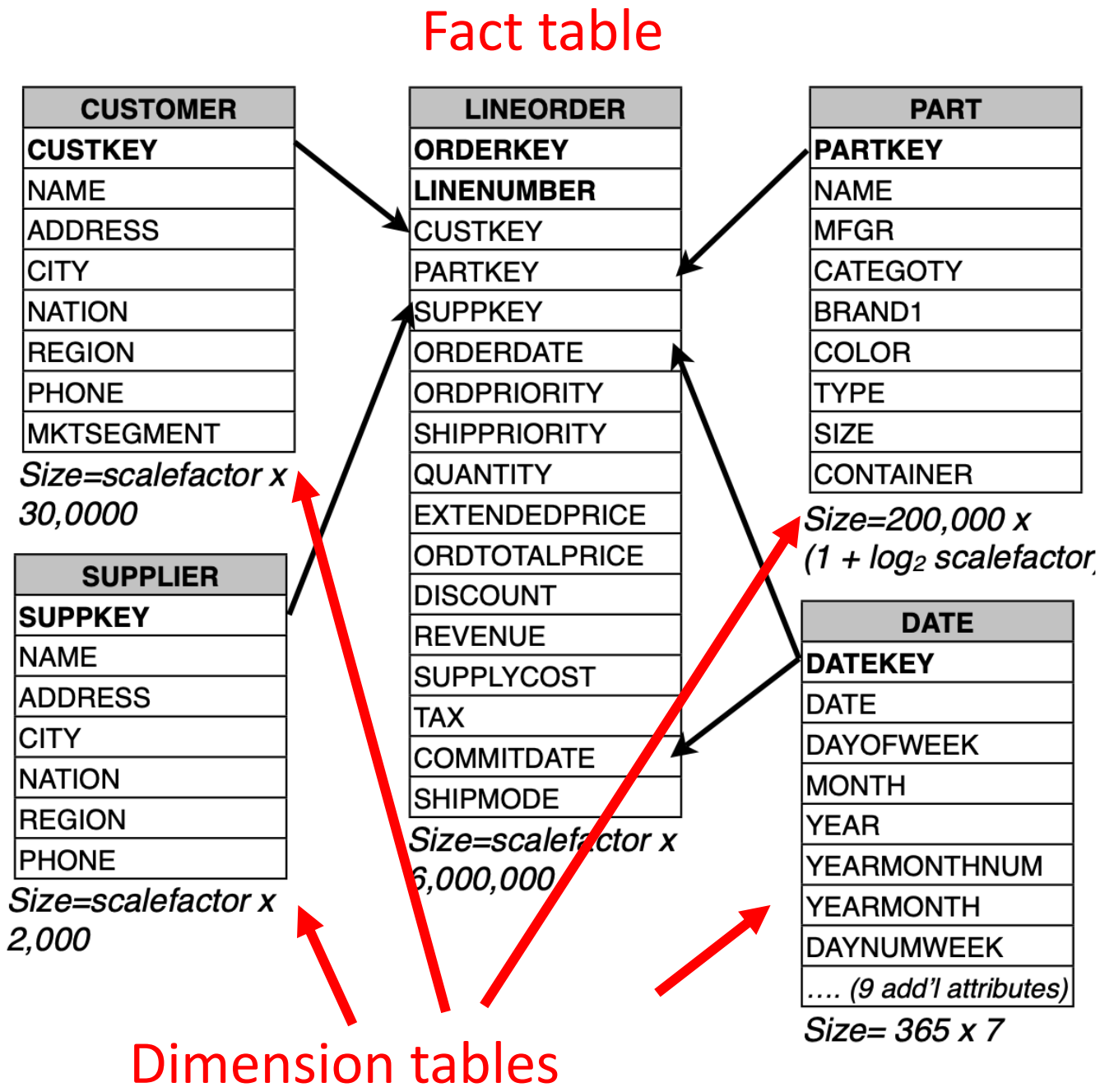
Star Schema Benchmark

Star-Schema Benchmark

13 queries

```
select sum(lo_extendedprice*lo_discount) as revenue
from lineorder, date
where lo_orderdate = d_datekey and
      d_year = 1993 and
      lo_discount between 1 and 3 and
      lo_quantity < 25;
```

```
select sum(lo_revenue), d_year, p_brand1
from lineorder, date, part, supplier
where lo_orderdate = d_datekey and
      lo_partkey = p_partkey and
      lo_suppkey = s_suppkey and
      p_category = 'MFGR#12' and
      s_region = 'AMERICA'
group by d_year, p_brand1
order by d_year, p_brand1;
```



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Invisible Joins

Idea: rewrite joins as predicates on foreign keys in fact table

Algorithm:

1. apply each predicate to the appropriate dimension table
2. build a hash table on matching keys
3. compute bitvector with bits set for qualifying positions (tuples)
4. intersect bitvectors (positions) via bitwise AND
5. for each resulting position reconstruct the resulting tuple

1. apply each predicate to the appropriate dimension table
2. build a hash table on matching keys

Apply `region = 'Asia'` on Customer table

custkey	region	nation	...
1	Asia	China	...
2	Europe	France	...
3	Asia	India	...

→ Hash table with keys 1 and 3

Apply `region = 'Asia'` on Supplier table

suppkey	region	nation	...
1	Asia	Russia	...
2	Europe	Spain	...

→ Hash table with key 1

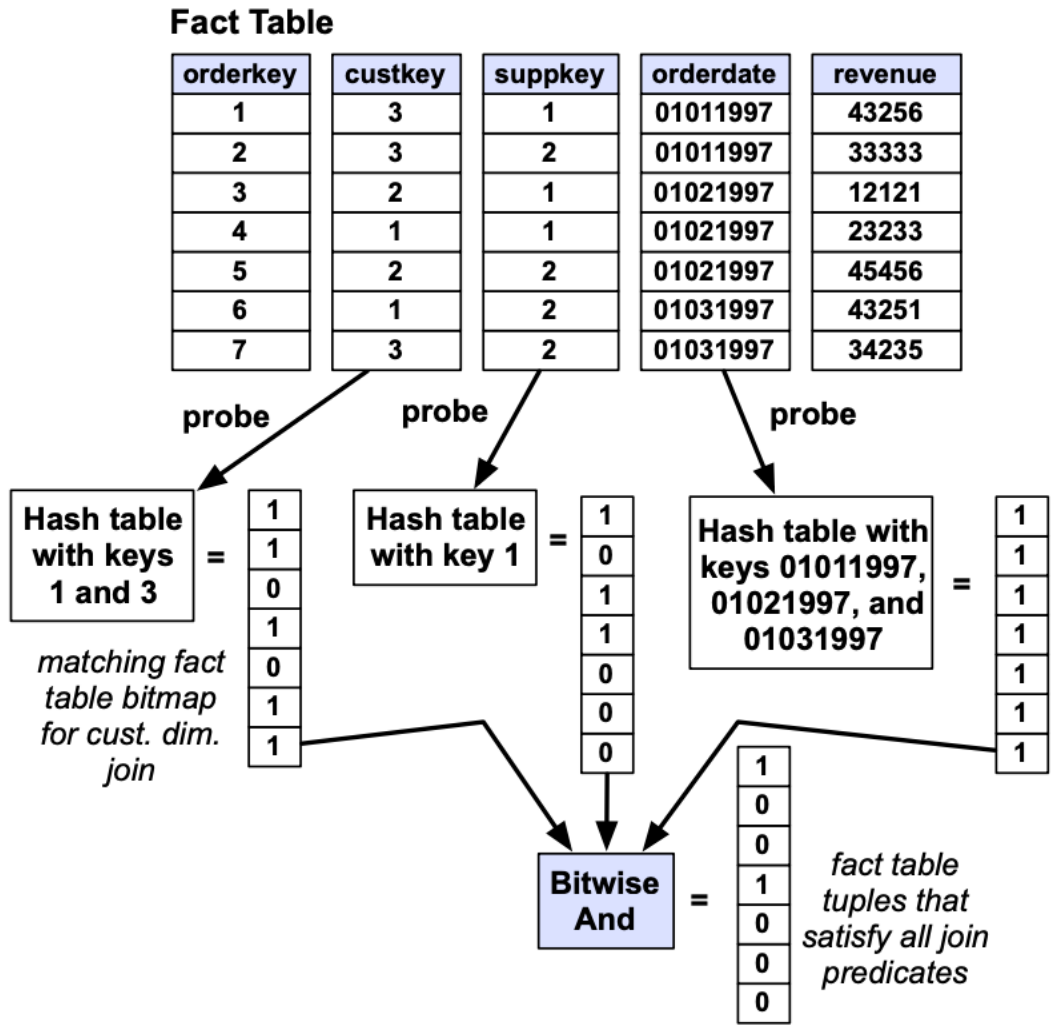
Apply `year in [1992,1997]` on Date table

dateid	year	...
01011997	1997	...
01021997	1997	...
01031997	1997	...

→ Hash table with keys 01011997, 01021997, and 01031997

```
SELECT c.nation, s.nation, d.year,
       sum(lo.revenue) as revenue
FROM customer AS c, lineorder AS lo,
     supplier AS s, dwdate AS d
WHERE lo.custkey = c.custkey AND
      lo.suppkey = s.suppkey AND
      lo.orderdate = d.datekey AND
      c.region = 'ASIA' AND s.region = 'ASIA' AND
      d.year >= 1992 and d.year <= 1997
GROUP BY c.nation, s.nation, d.year
ORDER BY d.year asc, revenue desc;
```

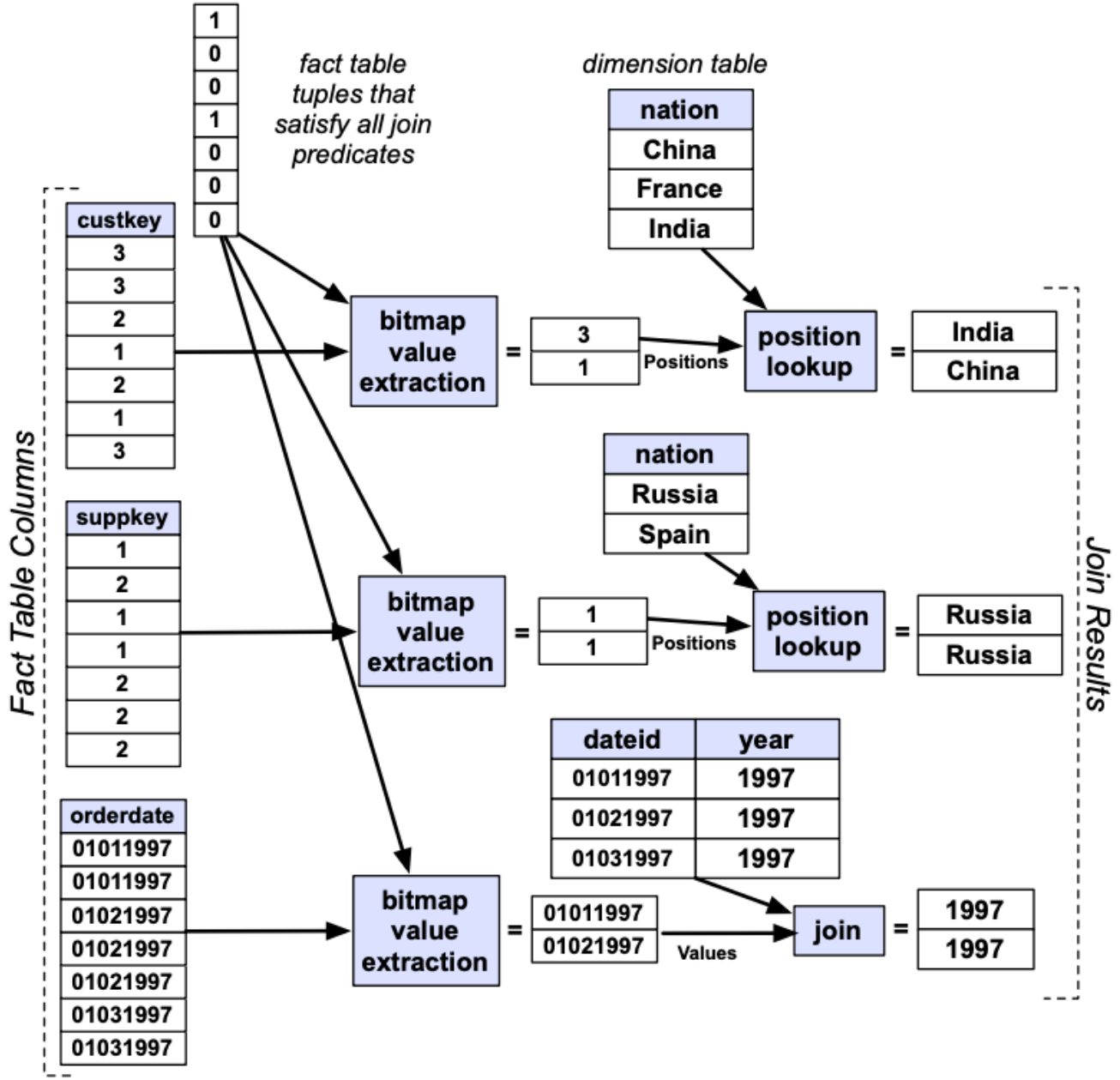
3. compute bitvector with bits set for qualifying positions (tuples)
4. intersect bitvectors (positions) via bitwise AND



5. For each resulting position, extract the values from the columns that are in the result

Are invisible joins a general join algorithm?

No! It works only for Star Schemas



Experiments

1 CPU 2.8GHz, 3GB RAM, Red Hat Linux 5

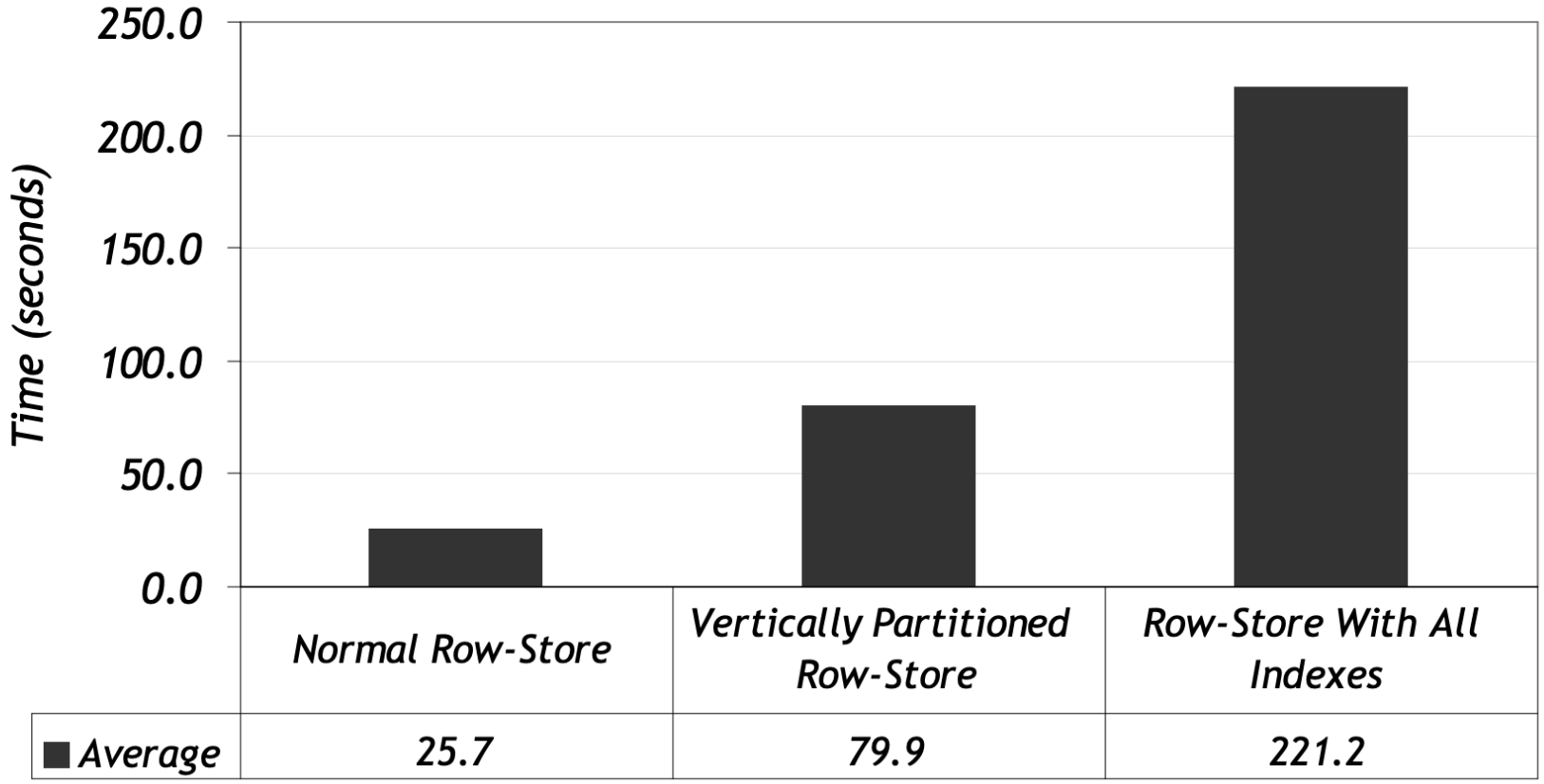
4-disk HDD array with 160-200MB/s aggregate bandwidth

(older paper, so small numbers!)

Report averages with “warm” bufferpool (smaller than data size)

Focus on SSB averages (the paper has more detailed graphs)

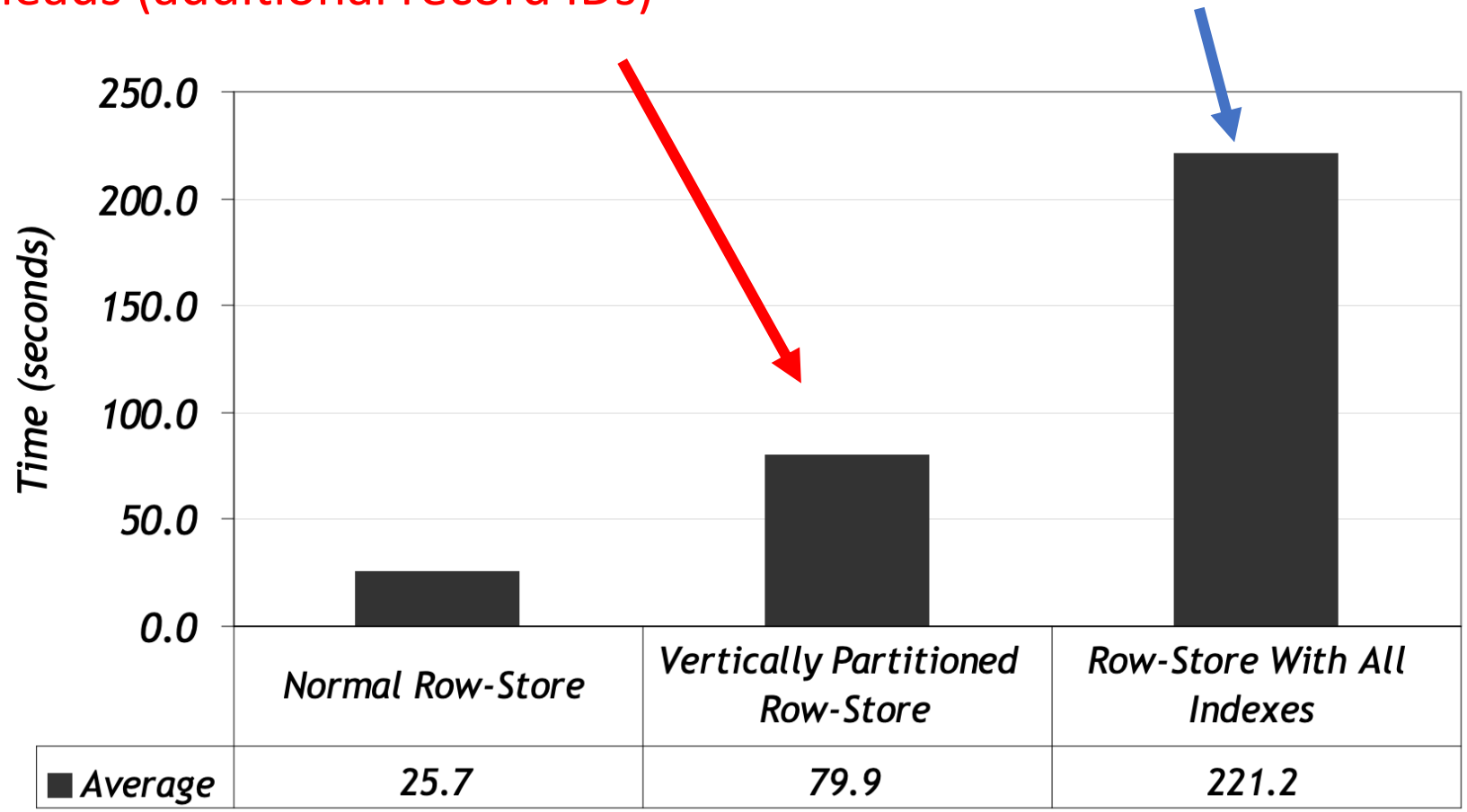
Experimenting with row-stores (SSB averages)



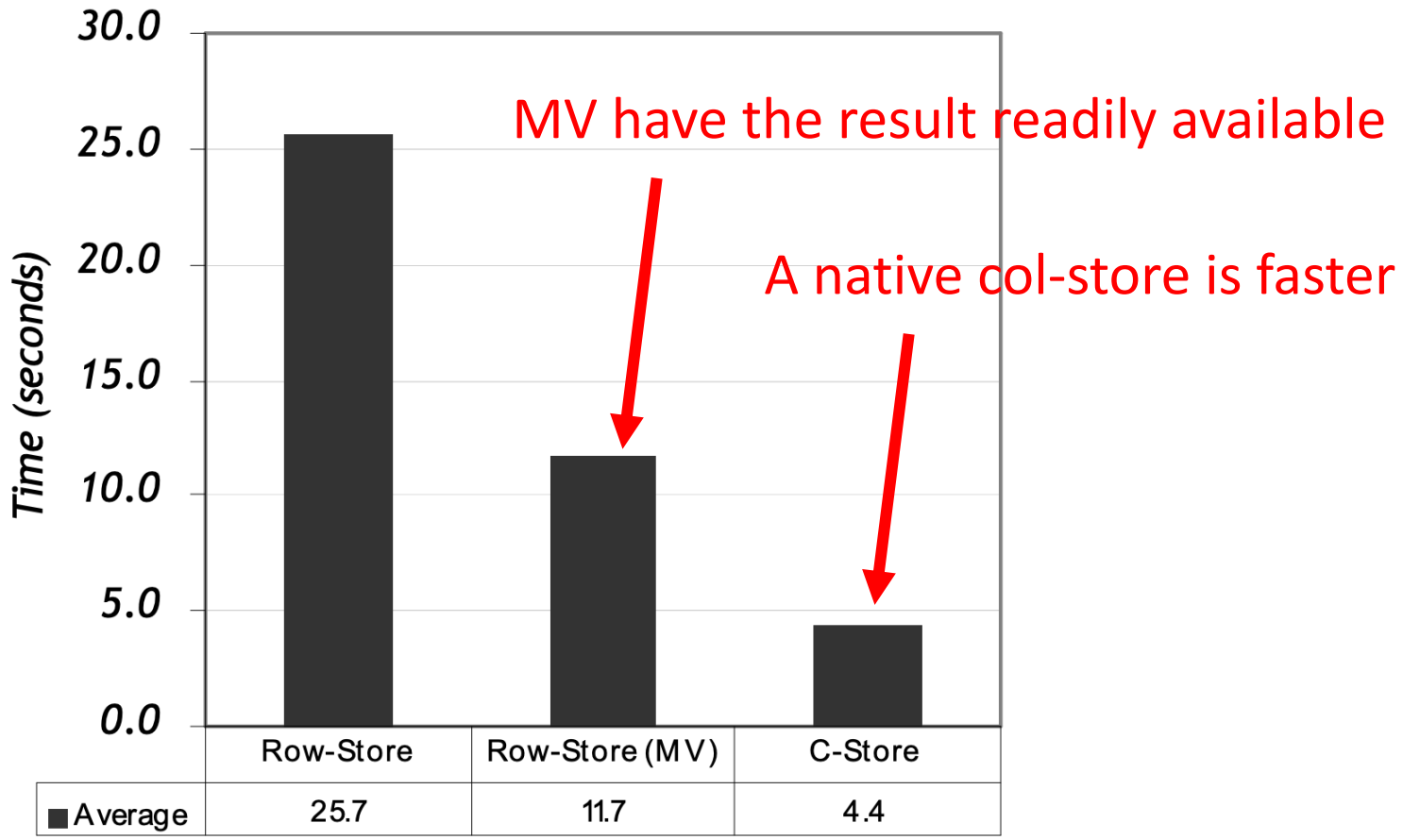
Experimenting with row-stores (SSB averages)

tuple reconstruction (via expensive joins)
prior to the join between tables

tuple overheads (additional record IDs)



Row-Stores vs. Column-Stores (SSB average)



Can we simulate a column-store with a row-store?

(a) All Indexes is a poor way to do it

(b) Vertical Partitioning's problem are NOT fundamental

i. *tuple header can be removed*
ii. *TIDs can be virtual*
iii. *horizontal partitioning can be based on the values of a different VP*

But still, column-stores and row-stores are apples and oranges!!

Methodology

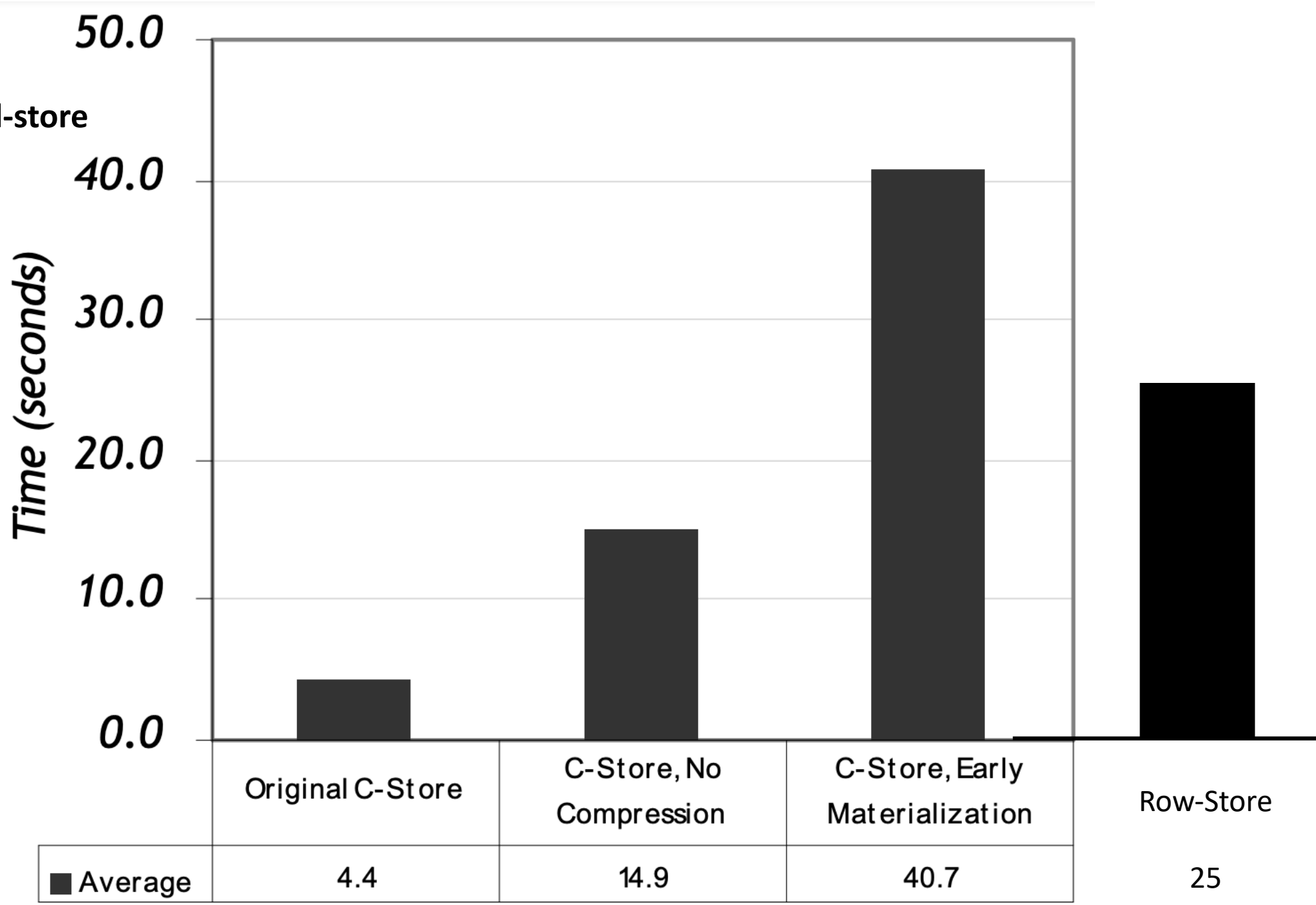
Start from a native column-store

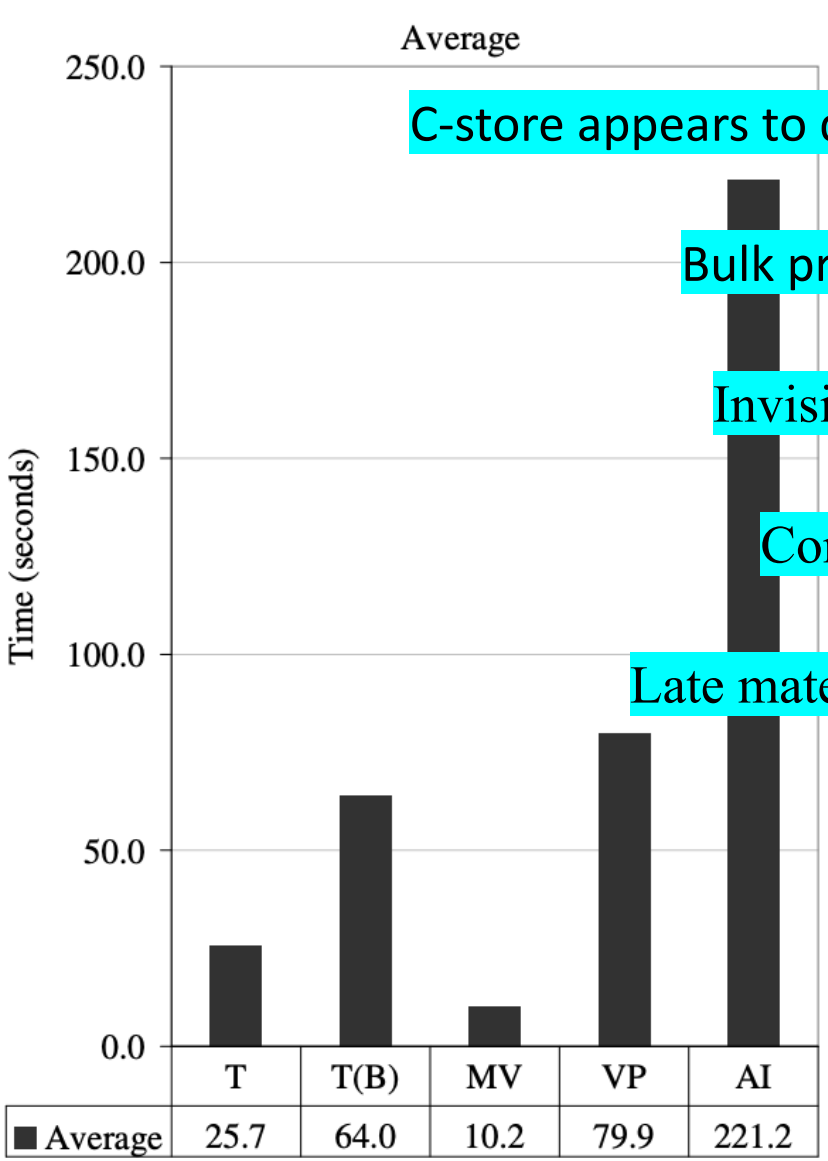
Remove column-store-specific performance optimizations

End with a column-store with a row-oriented query engine

To make the most of a col-store

1. Efficient support of vertical partitioning (compression)
2. Column-specific execution (late materialization)





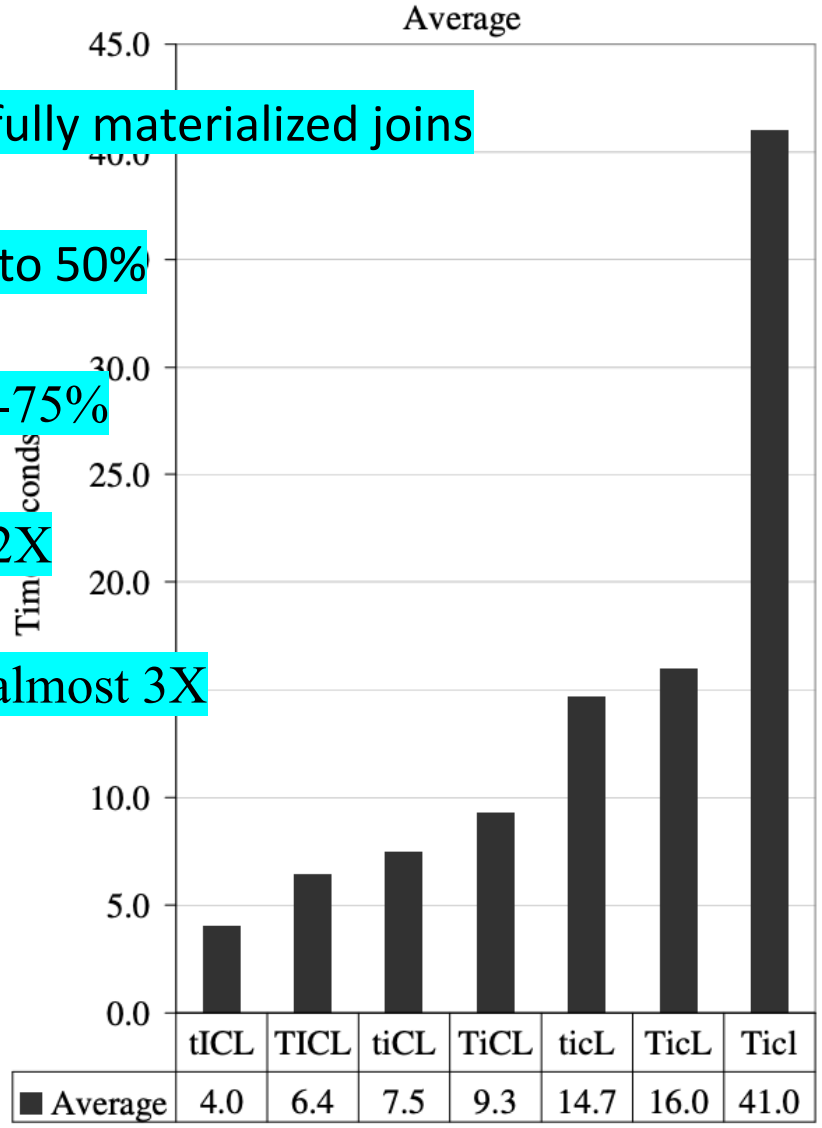
C-store appears to do even better than fully materialized joins

Bulk processing buys you 5 to 50%

Invisible join buys you 50-75%

Compression buys you 2X

Late materialization gets you almost 3X



T is traditional, T(B) is traditional (bitmap),
MV is materialized views, VP is vertical partitioning,
and AI is all indexes

T=tuple-at-a-time processing, t=block processing;
I=invisible join enabled, i=disabled;
C=compression enabled, c=disabled;
L=late materialization enabled, l=disabled

Things to remember

Row-stores vs. Col-stores: fundamental differences

✓Compression

✓ Late Materialization

✓Block Iteration

✓Column-store-specific join optimizations

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CS 561: Data Systems Architectures

class 5

Row-stores vs. Column-stores

Dr. Subhadeep Sarkar

https://bu-disc.github.io/CS561/

with slides based on Dan Abadi's