Bufferpool Implementation

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Introduction

• Bufferpool overview

- Buffer Hit
- Buffer Miss
- Dirty Bits
- Eviction Strategies
- Optimize for different workloads to maximize read and write performance by maximizing hits



LRU Policy

• When the bufferpool becomes full, we discard the least recently used page

 The pages that have been most heavily used in the past are more likely to be used heavily in the future too



- Deque -> easy insertion/remove from both end
- Store: Page Id
- Return: Position of the least used page in bufferpool -> target evicting position

CFLRU Policy

- Separate LRU list into 2 parts
- Keep a certain amount of dirty pages in cache to reduce the number of flash write operations
- we opted for the window to be ¹/₃
 of the current amount of pages in the buffer

LRU: p8 -> p7 -> p6-> p5 CFLRU: p7 -> p5 -> p8 -> p6



LRU-WSR Policy



- Second chance algorithm -> Cold flag can only be set on the second time (reordered writing sequences)
- The only difference between those two policies is that LRU-WSR assign each page with a bit flag called "cold-flag"
- In bufferpool we store tuple(page id, dirty bit, cold flag)



FIFO Policy

- First in first out, implemented by queue
- when a page hits:
 - LRU will move this page to the MRU position
 - FIFO will make no changes to the queue.
- Simple but not efficient for large number of pages → the operating system keeps track of all pages in the memory in a queue

1	2	3	4	5	1	3	1	6	3	2	3
							-0				

м	м	м	м	м	м	н	н	м	м	м	н
			4	4	4	4	4	4	3	3	3
		3	3	3	3	3	3	6	6	6	6
	2	2	2	2	1	1	1	1	1	1	1
1	1	1	1	5	5	5	5	5	5	2	2

M = Miss H = Hit

Disk Implementation

- Disk populated with arbitrary random bytes
- Writes and reads with seekg
 - One continuous character when written over(except last character is '\n')
 - One page per line, no sectors because just a quick computation and isn't productive to objectives
- Keeping file open may not be true reflection of disk access unless disk size exceeds RAM
 - We saw only a slight increase in timings
- Disk object belongs to Buffer class

498	заавававававававававававававававававава
499	///////////////////////////////////////
500	;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;
501	:puvfdsaiabbplclamiexxcldhcrrxommiktncvyewzvidgkpqroot
502	jzvudylshvvqgpiegbayzewdkvjcdrllrihwisqpnngwdqclrelsiiv
503	<pre>wcyuurewcrjnsoqxkgkygpucflnjjvnhxneueksibbytqoscvccbrv</pre>
504	/wnaowqcmghafzircxrdyyovskzuxsdvqtveroidupfbqnuvnoylmm;
505	_delkiijdujkjnysgtqhtudykqrxzrlhwqsgacrfyaqhpqbxjsecmka
506	<pre>wiyysxtytlyqnzkqvupgbkrmmxpsistsasssrmskzqaprmfnhwvkgr</pre>
507	eeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeeee
508	<pre>/eenmtybfsjnonziclczgprzcpylrckmgqavlayrthejwfryrvzzmq;</pre>
509	<pre>wedaabflvydvicgwmqvetospopvlnenjlqlltqwrqbmydsvrlswejc</pre>
510	<pre>>>dahbpzambqvcvpgawnrxiuypbvnixxxazebrfcfjualpprsngjlod;</pre>
511	anwrwdpjoboirholmrjlidcurwsjuqcwdaocdfnuhbcyiqlukwhuzlo
512	<pre>waszswaultfllhgebnwlxyqhwczbwksukmtcitytpehclngocccbcs</pre>
513	;\$
514	.zlmsdyvqiprmhiitnrxsgnhtfqrrfefgqryvpvlzkdlsntncnkuvzc
515	<pre>sisiupyvrjzgawroowmeurqzovoylsvfapnufmswxrezqvngtbkpva</pre>
516	<pre>>tmcgyfmojkmapzpamhdnmzorvgoizcwuqbapiofrartqqjscqxqcxq</pre>
517	.hutslvubaedujuimtmkynqcuytpocyyjtrbgoyirellohtaahmawc@
518	/hgddzueaafwirlyfluvwasgpgzufykdgqgjsbosbvploajtnfqjfls
519	<pre>wrwbgfeweelapfjqdhtydquqgythckeaddbjighmmsodzxwfgpdjh;</pre>
520	3aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa
521	<pre>mpwgqjkhimxrmezxhknogdlmansxveohvkpnuzueosvcywzigowort</pre>
522	<pre>icymkqofwpfbzfllompatrkljmwcbzknbiblbpszhacgfprvcgvvxh;</pre>
523	<pre>icuxeppvoavoygcymzwluepoivzlhhcplzmqodleeivcpxddybqsht</pre>
524	/bmrtjgsrypszjsuugihblpritxuczkaaxuvicpzagsbpkylsgtvrkr
525	<pre>/hstszcwdjhqnibhfdlctqulfbkekzpgijcbjgxoqeefohovmbzfsvc</pre>
526	<pre>'yhlaxejcyzaekhxzgmdtombelxmlvkevrqyqwhtxitbuaatjmxcaj</pre>
527	- il arrrionttal klal czachkhwwmlaowyhwawrahftuyfoiuchwzyfly

Experiments





Challenges and Lessons Learned

- Disk management
 - At first only appending seemed possible without a complete overwrite
- Seg faults in overall simulation run script
- Construct the bufferpool to be more adaptive
- Learned that it's much easier to have a small workload and know exactly how it should perform
- Read other research for guidance

Conclusions

- Implemented: FIFO, LRU, CFLRU, LRU-WSR, and Disk Functionality
- For balanced workloads, the simplest approach is also the best
- Intended extensions of this work:
 - Experiment with read heavy and write heavy workloads
 - Add more flexibility to the buffer with more parameters

100,000330901	
namespace b	pufmanager {
You, 6 days	s ago 1 author (You)
class E	Buffer {
	This class maintains specific property of the buffer.
	You definitely need to modify this part
	You need to add more variables here for your implementation. For example, current
pri	
	Buffer(Simulation Environment* _env);
	<pre>static Buffer* buffer_instance;</pre>
	olic:
	//initizate bufferpool with <page bit="" dirty="" page="" size,=""></page>
	<pre>vector< pair<int, bool=""> > bufferpool;</int,></pre>
	<pre>//initizate bufferpool for LRU WSR with <page bit,="" cold="" dirty="" flag="" page="" size,=""></page></pre>
	<pre>vector< tuple<int, bool="" bool,=""> > bufferpool_wsr;</int,></pre>
	//deck as the lru candidate list
	<pre>deque<int> lru_candidate;</int></pre>
	<pre>//for cflru: going to store either pageid of clean or -1</pre>
	deque <int> fifo_candidates;</int>
	<pre>static long max_buffer_size; //in pages</pre>
	static Puffer* astPufferInstance/Simulation Environment* onu);
	// statistics to track
	static int huffer hit.
	static int buffer miss
	static int read in:
	static int write in:
	static char disk write char:
	<pre>static std::chrono::duration <double. milli=""> timing:</double.></pre>
	// Page replacement algorithm
	<pre>int LRU();</pre>
	<pre>int LRUWSR();</pre>
	<pre>int FIFO();</pre>
	<pre>int CFLRU();</pre>
	// Disk parameters
	std::fstream disk;
	static int pageSize;
	// Statistics collection
	static void writeResults().
	static int printRuffer():
	static int printStats():
}:	
,,	