CS460: Intro to Database Systems

Class 22: Concurrency Control

Instructor: Manos Athanassoulis

https://bu-disc.github.io/CS460/
Concurrency Control

Serializability

Readings: Chapter 17.1

Two phase locking

Lock management and deadlocks

Locking granularity

Tree locking

Phantoms and predicate locking
Review

DBMSs support **ACID Transaction semantics**

Concurrency control and Crash Recovery are key components

For **Isolation** property, **serial execution** of transactions is safe but slow

- Try to find schedules equivalent to serial execution
Formal Properties of Schedules

*Serial schedule:* Schedule that does not interleave the actions of different transactions

*Equivalent schedules:* For any database state, the effect of executing the first schedule is identical to the effect of executing the second schedule

*Serializable schedule:* A schedule that is equivalent to some serial execution of the transactions

Note: If each transaction preserves consistency, every serializable schedule preserves consistency.
Conflicting Operations

We need a formal notion of equivalence that can be implemented efficiently
– Base it on the notion of “conflicting” operations

Definition: Two operations conflict if:
– They are done by different transactions,
– They are done on the same object,
– And at least one of them is a write
Conflict Serializable Schedules

**Definition:** Two schedules are *conflict equivalent* iff:
- They involve the same actions of the same transactions, and
- every pair of conflicting actions is ordered the same way

**Definition:** Schedule S is *conflict serializable* if:
- S is conflict equivalent to some serial schedule

Note, some “serializable” schedules are NOT conflict serializable
- A price we pay to achieve efficient enforcement
Conflict Serializability – Intuition

A schedule S is conflict serializable if:

– You are able to transform S into a serial schedule by swapping *consecutive non-conflicting* operations of different transactions

Example:

\[
\begin{align*}
R(A) & \quad W(A) & R(B) & \quad W(B) \\
R(A) & \quad W(A) & R(B) & \quad W(B) \\
R(A) & \quad W(A) & R(B) & \quad W(B) \\
\equiv & & & \\
R(A) & \quad W(A) & R(B) & \quad W(B) \\
\end{align*}
\]
Conflict Serializability (Continued)

Here’s another example:

 Serializable or not?

 NOT!
Dependency Graph

**Dependency graph:**
- One node per transaction
- Edge from $T_i$ to $T_j$ if:
  - An operation $O_i$ of $T_i$ conflicts with an operation $O_j$ of $T_j$ and
  - $O_i$ appears earlier in the schedule than $O_j$

**Theorem:** Schedule is conflict serializable if and only if its dependency graph is acyclic
Example

A schedule that is not conflict serializable:

| T1:       | R(A), W(A), R(B), W(B) |
| T2:       | R(A), W(A), R(B), W(B) |

The cycle in the graph reveals the problem. The output of T1 depends on T2, and vice-versa.
View Serializability

Alternative (weaker) notion of serializability

Schedules S1 and S2 are view equivalent if:

1. If $T_i$ reads initial value of A in $S_1$, then $T_i$ also reads initial value of A in $S_2$
2. If $T_i$ reads value of A written by $T_j$ in $S_1$, then $T_i$ also reads value of A written by $T_j$ in $S_2$
3. If $T_i$ writes final value of A in $S_1$, then $T_i$ also writes final value of A in $S_2$

Basically, allows all conflict serializable schedules + “blind writes”

| T1: R(A) | W(A) | view | T1: R(A),W(A) |
| T2: W(A) | | | T2: W(A) |
| T3: W(A) | | | T3: W(A) |
Notes on Serializability Definitions

**View Serializability** allows (slightly) more schedules than **Conflict Serializability**

- Problem: it is difficult to enforce efficiently

Neither definition allows all schedules that you would consider “serializable”

- Because they don’t understand the meanings of the operations or the data

In practice, **Conflict Serializability** is used, because it can be enforced **efficiently**

- To allow more concurrency, some special cases do get handled separately, such as travel reservations
Serializability Summary

Serial schedule

Equivalent schedules

Conflict Serializable schedule $S_a$, if $S_a$ conflict equivalent with (some) $S_{serial}$

View Serializable schedule $S_b$, if $S_b$ view equivalent with (some) $S_{serial}$

easier to enforce!

Conflict equivalent = if all conflicting op’s same order
View equivalent = if same view after
Concurrent Control

Serializability

Two phase locking

Readings: Chapter 17.1

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Locking granularity

Tree locking

Phantoms and predicate locking
Two-Phase Locking (2PL)

Locking Protocol

- each transaction obtains
  - S (*shared*) lock on object before reading
  - X (*exclusive*) lock on object before writing

- A transaction cannot request additional locks once it releases any locks
- Thus, there is a “growing phase” followed by a “shrinking phase”

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<thead>
<tr>
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<tbody>
<tr>
<td>S</td>
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Two-Phase Locking (2PL)

2PL on its own is sufficient to guarantee conflict serializability (i.e., schedules whose dependency graph is acyclic), but, it is subject to Cascading Aborts.
Strict 2PL

Problem: Cascading Aborts
Example: rollback of T1 requires rollback of T2!

| T1:  | R(A), W(A),          | R(B), W(B), Abort |
| T2:  | R(A), W(A)           |

How to avoid Cascading Aborts?

Strict Two-phase Locking (Strict 2PL) Protocol:
- Same as 2PL, except:
- All locks held by a transaction are released only when the transaction completes
Allows only conflict serializable schedules, but it is actually stronger than needed for that purpose

In effect, “shrinking phase” is delayed until

a) Transaction has committed (commit log record on disk), or
b) Decision has been made to abort the transaction (locks can be released after rollback)
**Non-2PL, A= 1000, B=2000, Output =?**

<table>
<thead>
<tr>
<th>Operation</th>
<th>Event</th>
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<tbody>
<tr>
<td>Lock_X(A)</td>
<td></td>
</tr>
<tr>
<td>Read(A)</td>
<td>Lock_S(A)</td>
</tr>
<tr>
<td>A: = A-50</td>
<td></td>
</tr>
<tr>
<td>Write(A)</td>
<td></td>
</tr>
<tr>
<td>Unlock(A)</td>
<td>Read(A)</td>
</tr>
<tr>
<td></td>
<td>Unlock(A)</td>
</tr>
<tr>
<td></td>
<td>Lock_S(B)</td>
</tr>
<tr>
<td>Lock_X(B)</td>
<td>Read(B)</td>
</tr>
<tr>
<td></td>
<td>Unlock(B)</td>
</tr>
<tr>
<td></td>
<td>PRINT(A+B)</td>
</tr>
<tr>
<td>Read(B)</td>
<td></td>
</tr>
<tr>
<td>B := B +50</td>
<td></td>
</tr>
<tr>
<td>Write(B)</td>
<td></td>
</tr>
<tr>
<td>Unlock(B)</td>
<td></td>
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</tbody>
</table>

**What is the problem here?**

A+B not executed in *Isolation*
**2PL, A= 1000, B=2000, Output = ?**

<p>| | |</p>
<table>
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<tr>
<td><strong>PRINT(A+B)</strong></td>
<td></td>
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**Questions:**
- what is the problem here?
- what if it aborts?

**Answer:**

Cascade Abort
**Strict 2PL, A= 1000, B=2000, Output =?**

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Venn Diagram for Schedules

- All Schedules
- View Serializable
- Conflict Serializable
- Serial

Avoid Cascading Abort
Q: Which schedules does Strict 2PL allow?
Two phase locking: Summary

Locks implement the notions of conflict directly

2PL has:
- Growing phase where locks are acquired and no lock is released
- Shrinking phase where locks are released and no lock is acquired

Strict 2PL requires all locks to be released at once, when transaction ends
Concurrency Control

Serializability

Two phase locking

Lock management and deadlocks

Locking granularity

Tree locking

Phantoms and predicate locking

Readings: Chapter 17.2-17.4
Lock Management

Lock and unlock requests handled by the Lock Manager

Lock Manager contains an entry for each currently held lock

Lock table entry:

- Pointer to list of transactions currently holding the lock
- Type of lock held (shared or exclusive)
- Pointer to queue of lock requests
Lock Management, continued

Basic operation: when lock request arrives see if any other transaction holds a conflicting lock

- If not, create an entry and grant the lock
- Else, put the requestor on the wait queue

Lock upgrade: transaction that holds a shared lock can be upgraded to hold an exclusive lock

Two-phase locking is simple enough, right?
**Example: Output = ?**

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What is the problem here? Deadlock
Deadlocks

**Deadlock**: Cycle of transactions waiting for locks to be released by each other

Two ways of dealing with deadlocks:
- Deadlock prevention
- Deadlock detection

Many systems just “punt” and use Timeouts
- What are the dangers with this approach?

forward progress
Deadlock Detection

Create a **waits-for graph**:

- Nodes are transactions
- Edge from $T_i$ to $T_j$ if $T_i$ is waiting for $T_j$ to release a lock

Periodically check for cycles in waits-for graph

![Diagram](image-url)

$T_i$ waits $T_j$ to release a lock

Important!! This is different than dependency graph!
Deadlock Detection (Continued)

Example:

T1: \( S(A), S(D), S(B) \)

T2: \( X(B) \)

T3: \( S(D), S(C), X(A) \)

T4: \( X(B) \)

waits for

waits for

waits for

waits for

deadlock!
Deadlock Prevention

Assign priorities based on timestamps

Say $T_i$ wants a lock that $T_j$ holds

Two policies are possible:

**Wait-Die:** If $T_i$ has higher priority, $T_i$ waits for $T_j$;
otherwise $T_i$ aborts

**Wound-wait:** If $T_i$ has higher priority, $T_j$ aborts;
otherwise $T_i$ waits

Why do these schemes guarantee no deadlocks?

Important detail: If a transaction re-starts, make sure it gets its original timestamp. -- Why?

to avoid starvation!
Deadlocks: summary

The lock manager keeps track of the locks issued

Deadlock is a cycle of transactions waiting for locks to be released to each other

Deadlocks may arise and can be:
  – Prevented, e.g. using timestamps
  – Detected, e.g. using waits-for graphs
CS460: Intro to Database Systems

Class 23: Concurrency Control (cont’d)

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Conflict equivalent = if all conflicting op’s same order

View equivalent = if same view after

easier to enforce!

From previous class!
Reminder: Two-Phase Locking (2PL)

Lock Compatibility Matrix

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<td>–</td>
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From previous class!
Reminder: Strict Two-Phase Locking (Strict-2PL)

- Acquire locks during the acquisition phase.
- Release all locks at the end of the transaction (xact).

Lock Compatibility Matrix:

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From previous class!
Reminder: Deadlock Detection (waits-for-graph)

Example:

T1: \( S(A), S(D), \)  
T2: \( X(B) \)  
T3: \( S(D), S(C), X(A) \)  
T4: \( X(B) \)  

From previous class!
Reminder: Deadlock Prevention

Assign *priorities* based on *timestamps*

Say Ti wants a lock that Tj holds

Two policies are possible:

**Wait-Die**: If Ti has higher priority, Ti waits for Tj; otherwise Ti **aborts**

**Wound-wait**: If Ti has higher priority, Tj aborts; otherwise Ti **waits**

the *trx* that wants a (held) lock:

high priority : waits

low priority : aborts

high priority : kills the other

low priority : wait

From previous class!
Concurrency Control

Serializability

Two phase locking

Lock management and deadlocks

Locking granularity

Readings: Chapter 17.5.2

Tree locking

Phantoms and predicate locking
Multiple-Granularity Locks

Hard to decide what granularity to lock (tuples vs. pages vs. tables)

Shouldn’t have to make same decision for all transactions!

Data “containers” are nested:

- Database
- Tables
- Pages
- Tuples

what if T1 has a lock on a page, and T2 on a tuple of this page? how to correctly protect the data?
Solution: New Lock Modes, Protocol

Allow transaction to lock at each level, but with a special protocol using new “intention” locks:
Still need S and X locks, but before locking an item, transaction must have proper intension locks on all its ancestors in the granularity hierarchy

IS – Intent to get S lock(s) at finer granularity
IX – Intent to get X lock(s) at finer granularity
SIX mode: Like S & IX at the same time.

Why is it useful?
Multiple Granularity Lock Protocol

Each transaction starts from the root of the hierarchy
To get S or IS lock on a node, must hold IS or IX on parent node
  – What if transaction holds SIX on parent? S on parent?
To get X or IX or SIX on a node, must hold IX or SIX on parent node
Must release locks in bottom-up order and must follow 2PL

Protocol is equivalent to directly setting locks at the leaf levels of the hierarchy.
### Lock Compatibility Matrix

<table>
<thead>
<tr>
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</table>

- **IS** – Intent to get S lock(s) at finer granularity
- **IX** – Intent to get X lock(s) at finer granularity
- **SIX mode**: S & IX at the same time
Examples – 2 level hierarchy

**T1** scans R, and updates a few tuples:
- T1 gets an SIX lock on R, then get X lock on tuples that are updated

**T2** uses an index to read only part of R:
- T2 gets an IS lock on R, and repeatedly gets an S lock on tuples of R

**T3** reads all of R:
- T3 gets an S lock on R
- OR, T3 could behave like T2
- We can use lock escalation to decide
- Lock escalation dynamically asks for coarser-grained locks when too many low level locks acquired

<table>
<thead>
<tr>
<th>Tuples</th>
<th>IS</th>
<th>IX</th>
<th>SIX</th>
<th>S</th>
<th>X</th>
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<tbody>
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<td>✓</td>
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Multiple granularity locking: Summary

Allows **flexibility** for each transaction to choose locking granularity independently

Introduces **hierarchy** of objects

Introduces **intention** locks
Concurrency Control

Serializability

Two phase locking

Lock management and deadlocks

Locking granularity

Tree locking

Readings: Chapter 17.5.2

Phantoms and predicate locking
Potential concurrency bottleneck:

Simply 2PL nodes while traversing:
- root node (and higher level nodes) become bottlenecks because every tree access begins at the root

How can we efficiently lock a particular leaf node?

Don’t confuse this with multiple granularity locking!!
Two Useful Observations

1. In a B+Tree, higher levels of the tree only direct searches for leaf pages
2. For inserts, a node on a path from root to modified leaf must be locked (in X mode, of course), only if a split can propagate up to it from the modified leaf (Similar point holds w.r.t. deletes)

We can exploit these observations to design efficient locking protocols that guarantee serializability **even though they violate 2PL**
A Simple Tree Locking Algorithm: “crabbing”

**Search:** Start at root and go down; repeatedly, S lock child then unlock parent

**Insert/Delete:** Start at root and go down, obtaining X locks as needed. Once child is locked, check if it is safe:

- If child is safe, release all locks on ancestors

**Safe node:** Node such that changes will not propagate up beyond this node.

-- When?

- Insertions: Node is not full
- Deletions: Node is not half-empty
Example

Do:
1) Search 38*
2) Delete 38*
3) Insert 45*
4) Insert 25*

shared lock
Example

Do:
1) Search 38*
2) Delete 38*
3) Insert 45*
4) Insert 25*

exclusive lock

safe?
Example

Do:
1) Search 38*
2) Delete 38*
3) Insert 45*
4) Insert 25*

safe?
Example

Do:
1) Search 38*
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4) Insert 25*

Do these operations under exclusive lock?
Example

Do:
1) Search 38*
2) Delete 38*
3) Insert 45*
4) Insert 25*

why does this scheme violate 2PL?
releases locks and acquires new locks!
Concurrency Control

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Readings: Chapter 17.5.1
Dynamic Databases – The “Phantom” Problem

If we relax the assumption that the DB is a fixed collection of objects, even Strict 2PL (on individual items) will not ensure serializability:

Consider T1 – “Find oldest sailor”

- T1 locks all records, and finds oldest sailor (say, age = 71)
- Next, T2 inserts a new sailor; age = 96 and commits
- T1 (within the same transaction) checks for the oldest sailor again and finds sailor aged 96!

The sailor with age 96 is a “phantom tuple” from T1’s point of view:

“first it’s not there then it is”

No serial execution of T1 and T2 could result to this!
The “Phantom” Problem – ex. 2

Consider T3 – “Find oldest sailor for each rating”

- T3 locks all pages containing sailor records with \textit{rating} = 1, and finds \textit{oldest} sailor (say, \textit{age} = 71)
- Next, T4 inserts a new sailor; \textit{rating} = 1, \textit{age} = 96
- T4 also deletes oldest sailor with \textit{rating} = 2 (and, say, \textit{age} = 80), and commits
- T3 now locks all pages containing sailor records with \textit{rating} = 2, and finds \textit{oldest} (say, \textit{age} = 63)

T3 saw only part of T4’s effects!

No serial execution where T3’s result could happen!
The Problem

T1 and T3 implicitly assumed that they had locked the set of all sailor records satisfying a predicate

– Assumption only holds if no sailor records are added while they are executing!
– Need some mechanism to enforce this assumption (Index locking and predicate locking)

Examples show that conflict serializability on reads and writes of individual items guarantees serializability only if the set of objects is fixed!
Predicate Locking

Predicate locking:
Grant lock on all records that satisfy some logical predicate, e.g., \( age > 2 \times salary \)

Index locking is a special case of predicate locking for which an index supports efficient implementation of the predicate lock

- What is the predicate in the sailor example?

In general, predicate locking has a lot of locking overhead
Index Locking

If there is a dense index on the \textit{rating} field using Alternative (2), T3 should lock the index page containing the data entries with \textit{rating} = 1

- If there are no records with \textit{rating} = 1, T3 must lock the index page where such a data entry \textit{would} be, if it existed!

If there is no suitable index, T3 must obtain:

1. A lock on every page in the table file
   \rightarrow prevent a record’s rating from being changed to 1

AND

2. The lock for the file itself
   \rightarrow prevent records with \textit{rating} = 1 from being added or deleted
Transaction Support in SQL-92

SERIALIZABLE
No phantoms, all reads repeatable, no “dirty” (uncommitted) reads

REPEATABLE READS
phantoms may happen

READ COMMITTED
phantoms and unrepeatable reads may happen

READ UNCOMMITTED
all of them may happen
Phantom problem: Summary

If database objects can be added/removed, need to guard against **Phantom Problem**

Must lock **logical** sets of records

Efficient solution: index locking